# Déjà Q: Using Dual Systems to Revisit q-Type Assumptions

Melissa Chase (MSR Redmond)

Sarah Meiklejohn (UC San Diego → University College London)

Historically, pairings have provided great functionality

Historically, pairings have provided great functionality

First IBE instantiation [BF01]

Historically, pairings have provided great functionality

- First IBE instantiation [BF01]
- Many other breakthroughs have followed [BBS04,GS08,KSW08,LW11,...]

Historically, pairings have provided great functionality

- First IBE instantiation [BF01]
- Many other breakthroughs have followed [BBS04,GS08,KSW08,LW11,...]

With great functionality, comes great (ir)responsibility!

Historically, pairings have provided great functionality

- First IBE instantiation [BF01]
- Many other breakthroughs have followed [BBS04,GS08,KSW08,LW11,...]

With great functionality, comes great (ir)responsibility!

First assumption: BDH (given (ga,gb,gc), compute e(g,g)abc)

Historically, pairings have provided great functionality

- First IBE instantiation [BF01]
- Many other breakthroughs have followed [BBS04,GS08,KSW08,LW11,...]

With great functionality, comes great (ir)responsibility!

- First assumption: BDH (given (ga,gb,gc), compute e(g,g)abc)
- Later assumptions: Subgroup Hiding [BGN05], Decision Linear, SXDH

Historically, pairings have provided great functionality

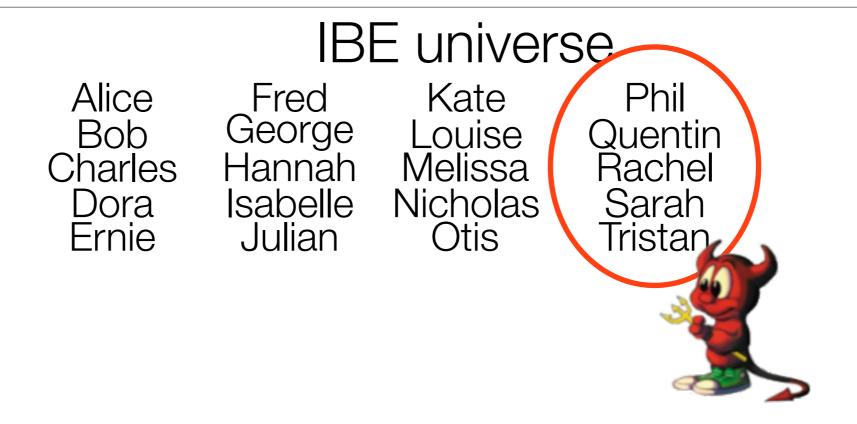
- First IBE instantiation [BF01]
- Many other breakthroughs have followed [BBS04,GS08,KSW08,LW11,...]

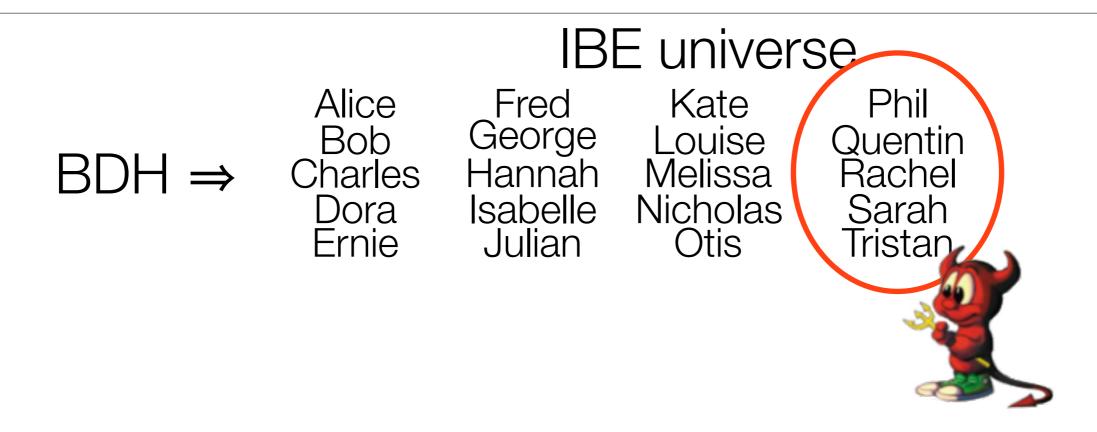
With great functionality, comes great (ir)responsibility!

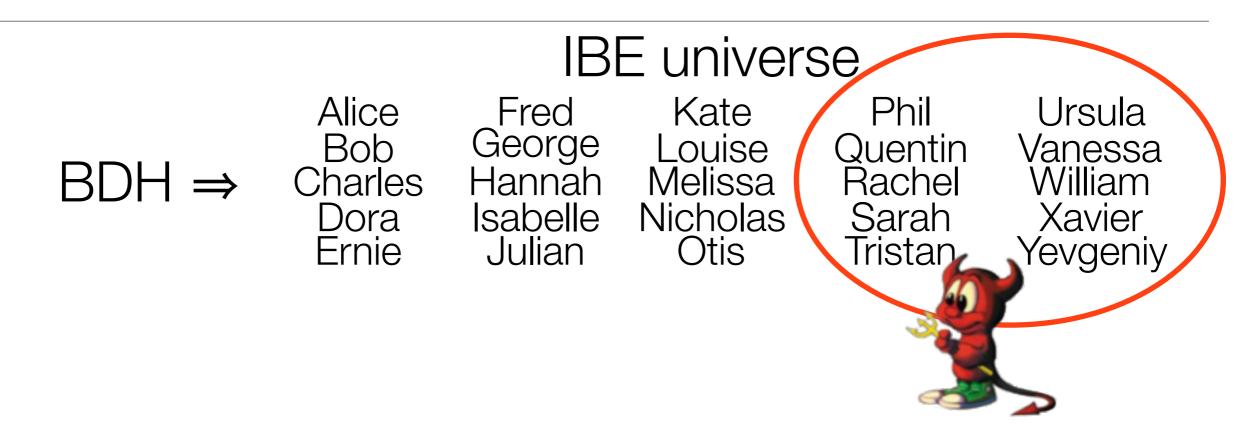
- First assumption: BDH (given (ga,gb,gc), compute e(g,g)abc)
- Later assumptions: Subgroup Hiding [BGN05], Decision Linear, SXDH
- Even later assumptions: q-SDH, q-ADHSDH, q-EDBDH, q-SDH-III, q-SFP, "source group q-parallel BDHE," etc.

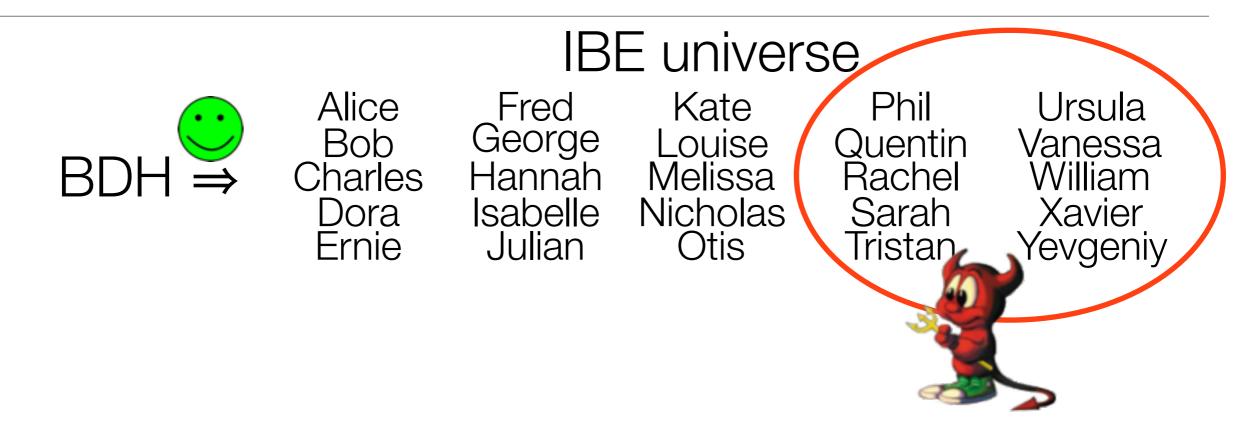
#### IBE universe

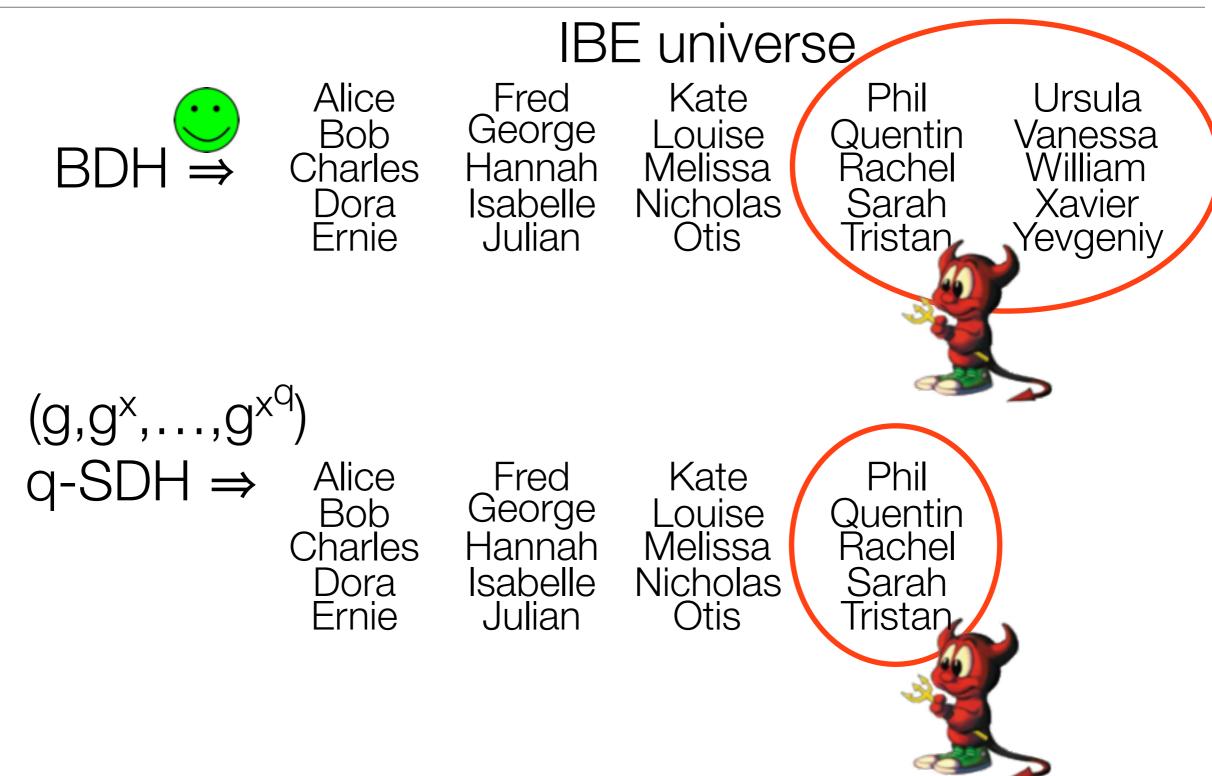
Alice Fred Kate Phil George Bob Louise Quentin Melissa Charles Hannah Rachel Nicholas Sarah Dora Isabelle Otis Tristan Ernie Julian

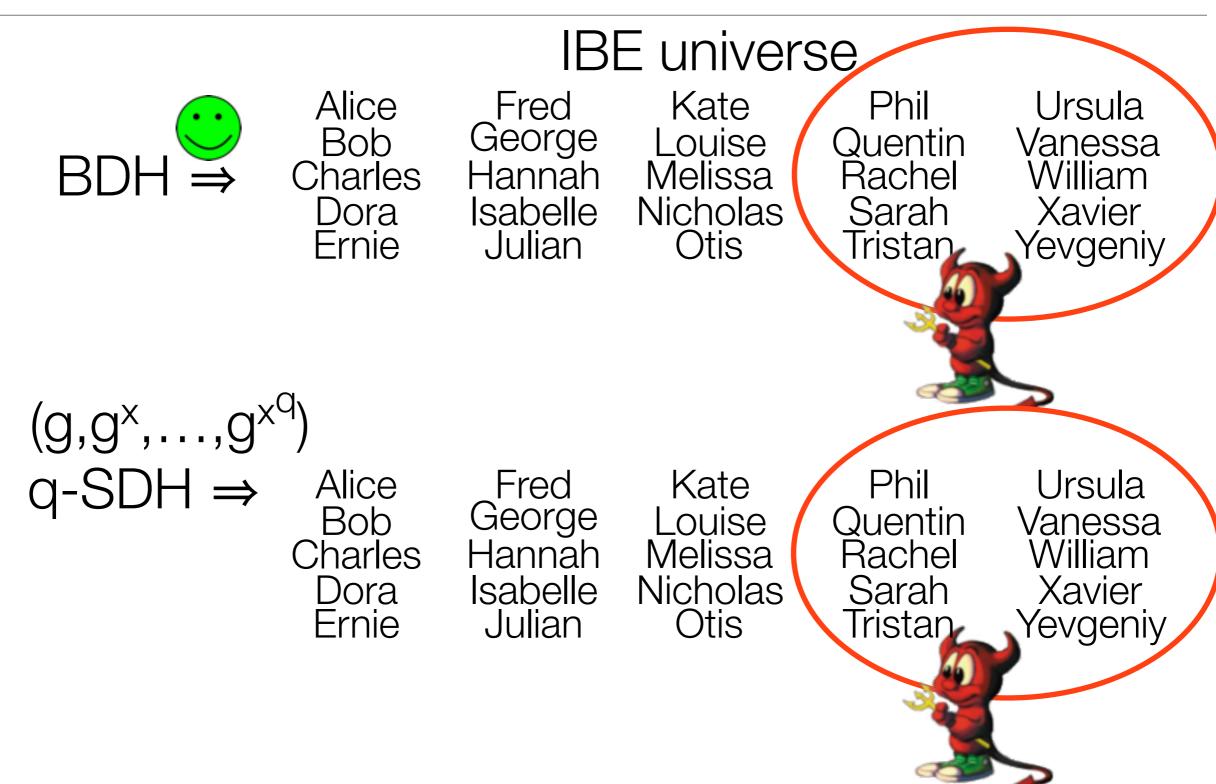


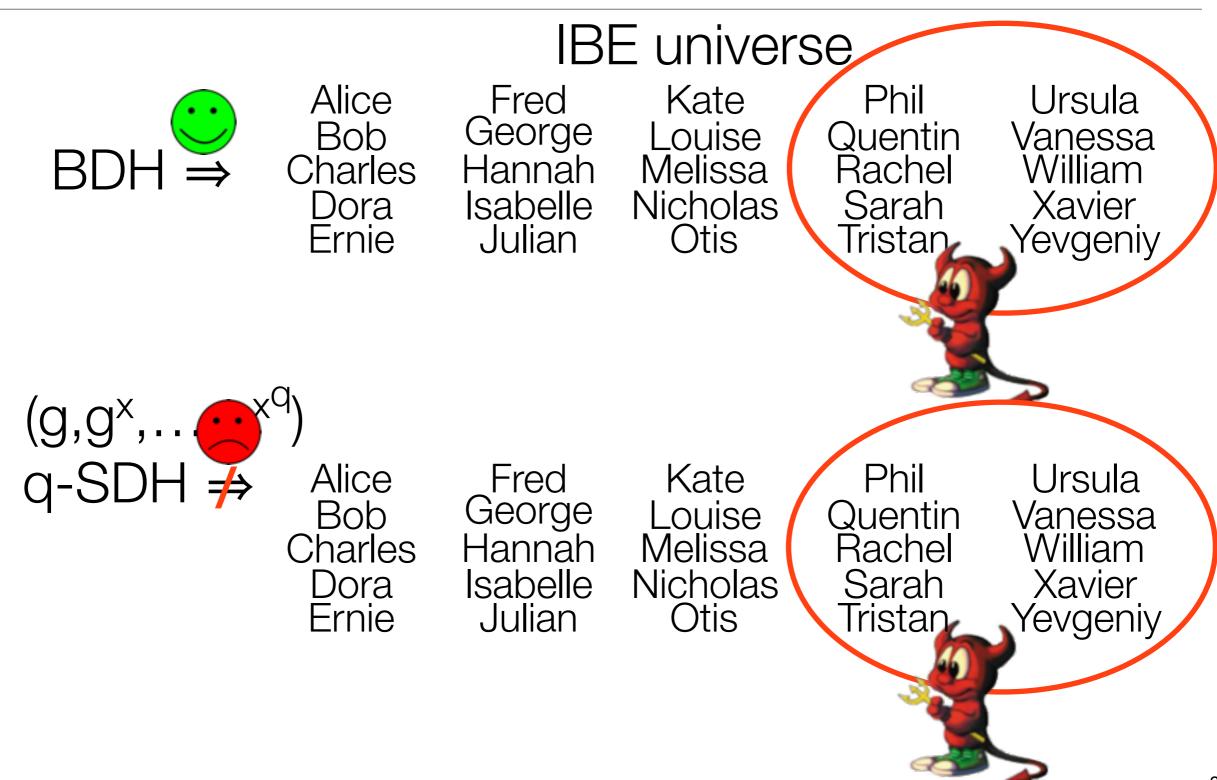


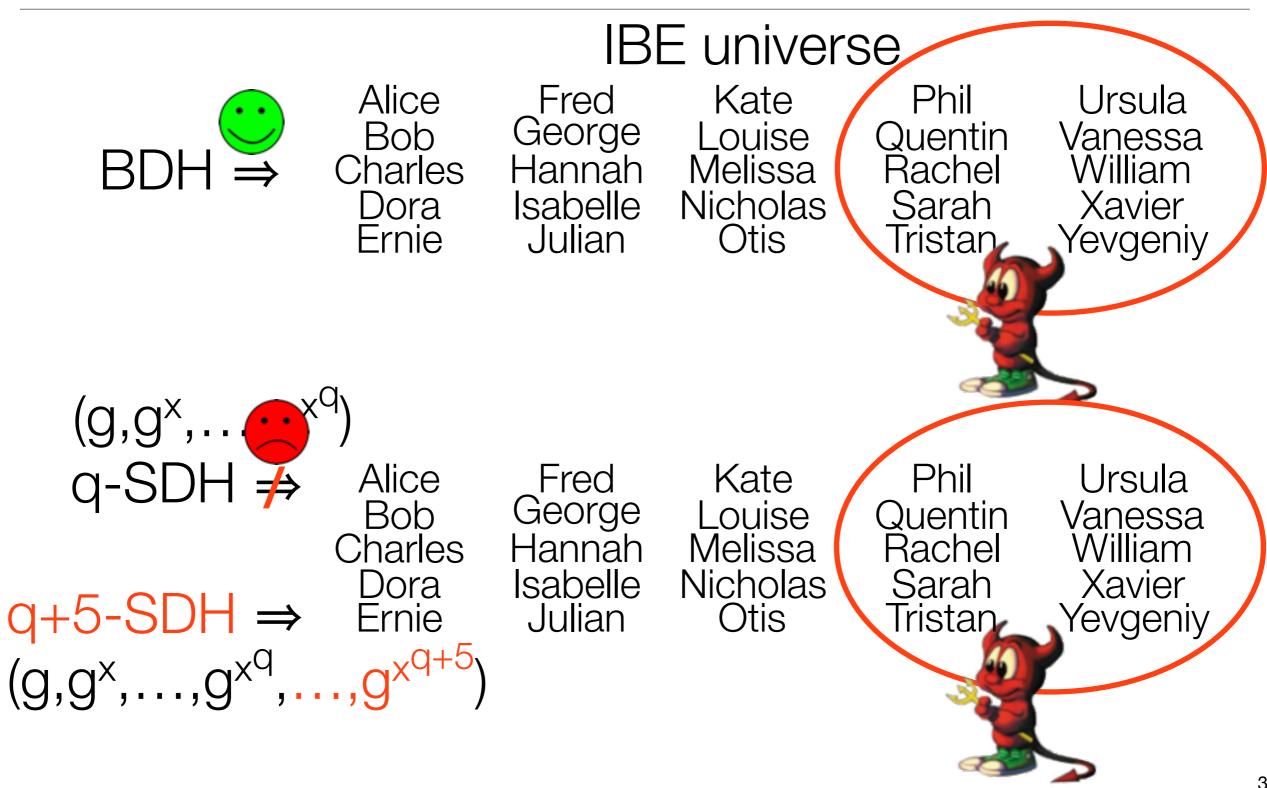


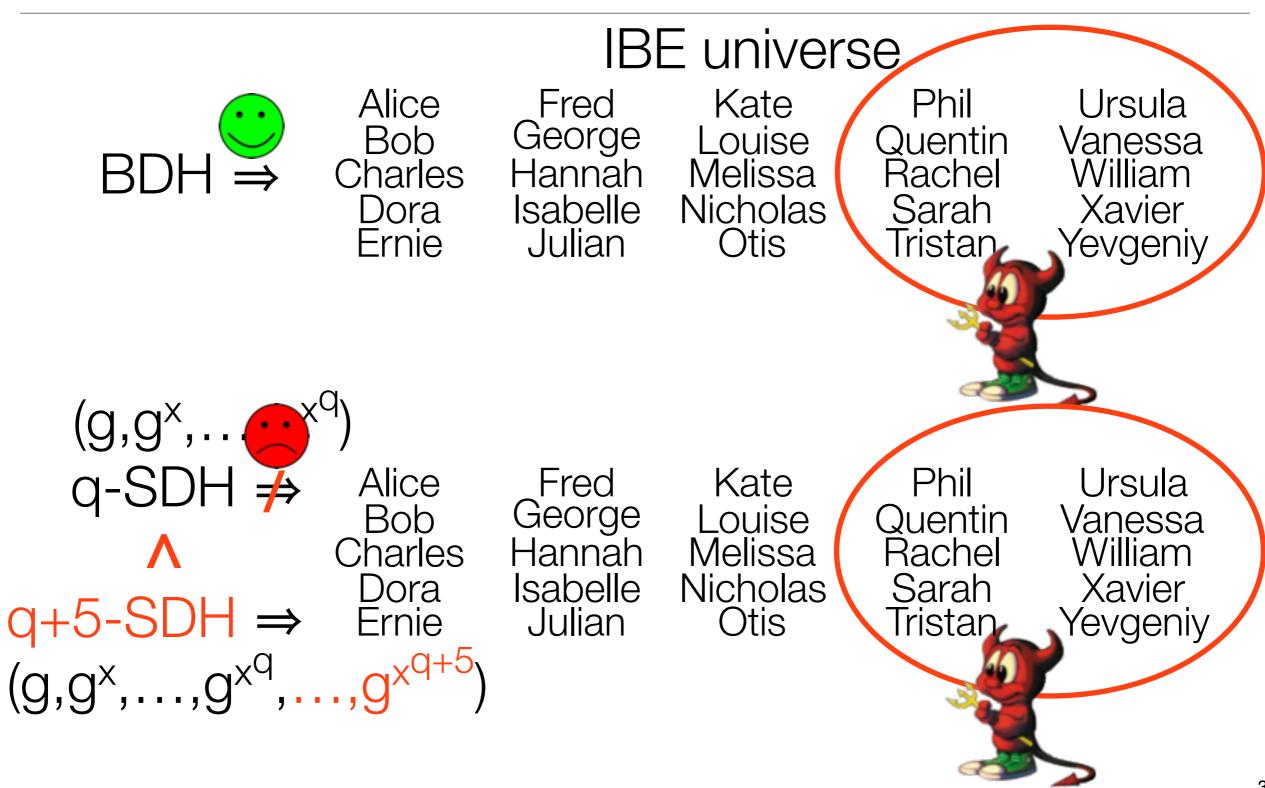


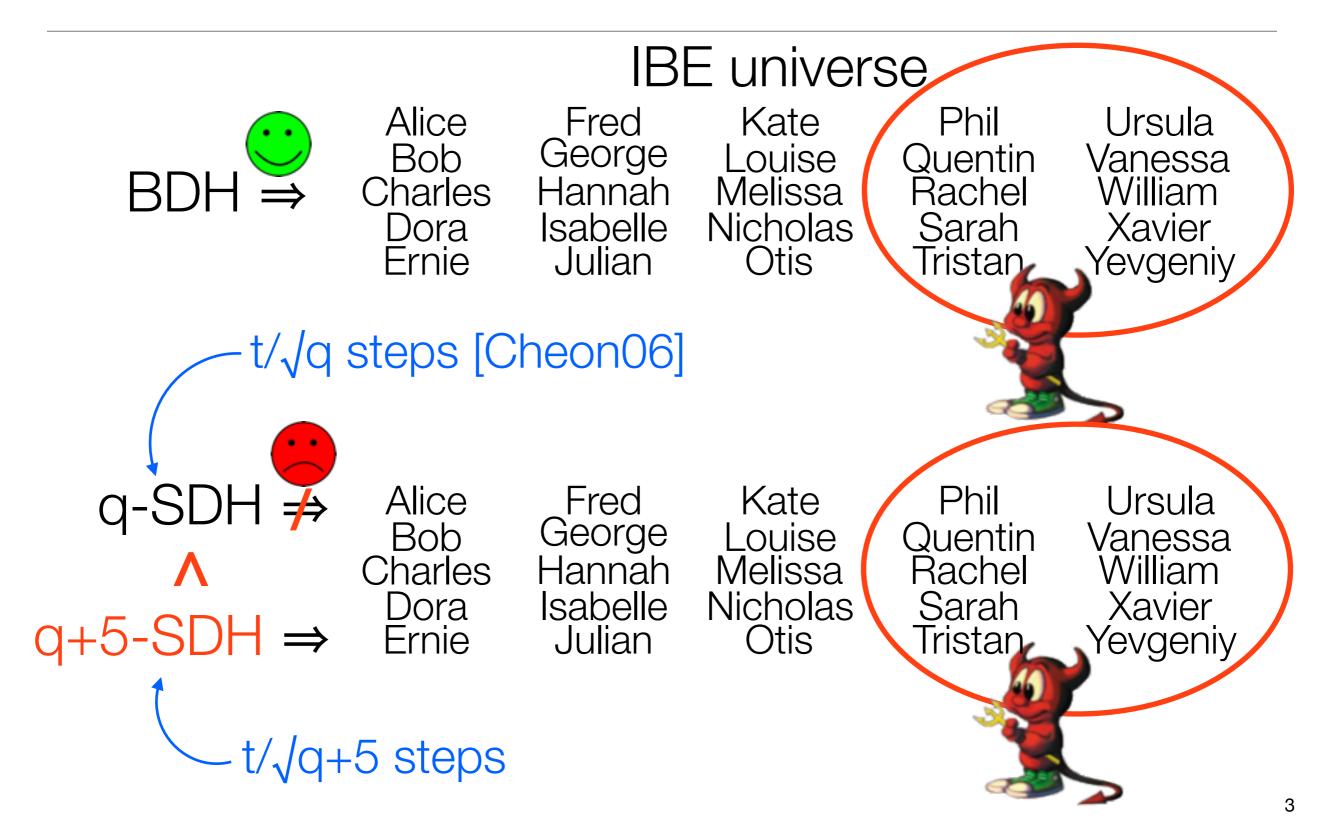












Dual systems [W09,...] have proved effective at removing q-type assumptions

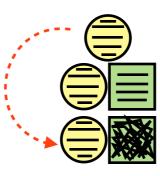
Dual systems [W09,...] have proved effective at removing q-type assumptions

Properties of bilinear groups: subgroup hiding and parameter hiding

Dual systems [W09,...] have proved effective at removing q-type assumptions

Properties of bilinear groups: subgroup hiding and parameter hiding

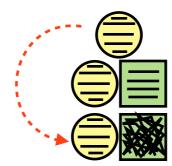
Abstract dual systems into three steps



Dual systems [W09,...] have proved effective at removing q-type assumptions

Properties of bilinear groups: subgroup hiding and parameter hiding

Abstract dual systems into three steps

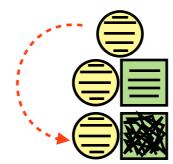


Apply dual systems directly to variants of the uber-assumption [BBG05,B08]

Dual systems [W09,...] have proved effective at removing q-type assumptions

Properties of bilinear groups: subgroup hiding and parameter hiding

Abstract dual systems into three steps



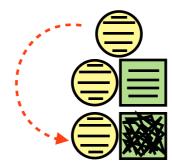
Apply dual systems directly to variants of the uber-assumption [BBG05,B08]

Reduce\* to an assumption that holds by a statistical argument

Dual systems [W09,...] have proved effective at removing q-type assumptions

Properties of bilinear groups: subgroup hiding and parameter hiding

Abstract dual systems into three steps



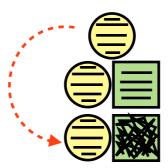
Apply dual systems directly to variants of the uber-assumption [BBG05,B08]

- Reduce\* to an assumption that holds by a statistical argument
- Adapt dual systems to work for deterministic primitives

Dual systems [W09,...] have proved effective at removing q-type assumptions

Properties of bilinear groups: subgroup hiding and parameter hiding

Abstract dual systems into three steps



Apply dual systems directly to variants of the uber-assumption [BBG05,B08]

- Reduce\* to an assumption that holds by a statistical argument
- Adapt dual systems to work for deterministic primitives

Extension to Dodis-Yampolskiy PRF [DY05]

Bilinear groups

Bilinear groups

q-Type assumptions

Bilinear groups

q-Type assumptions

Extensions

Bilinear groups q-Type assumptions Conclusions Extensions

Bilinear groups

Subgroup hiding Parameter hiding Dual systems

q-Type assumptions

Extensions

Conclusions

# Properties of (bilinear) groups

Standard bilinear group: (N, G, H, G<sub>T</sub>, e, g, h)

# Properties of (bilinear) groups

```
Standard bilinear group: (N, G, H, G<sub>T</sub>, e, g, h)

Group order;

prime or composite
```

#### Properties of (bilinear) groups

Standard bilinear group: (N, G, H, G<sub>T</sub>, e, g, h)

Group order;

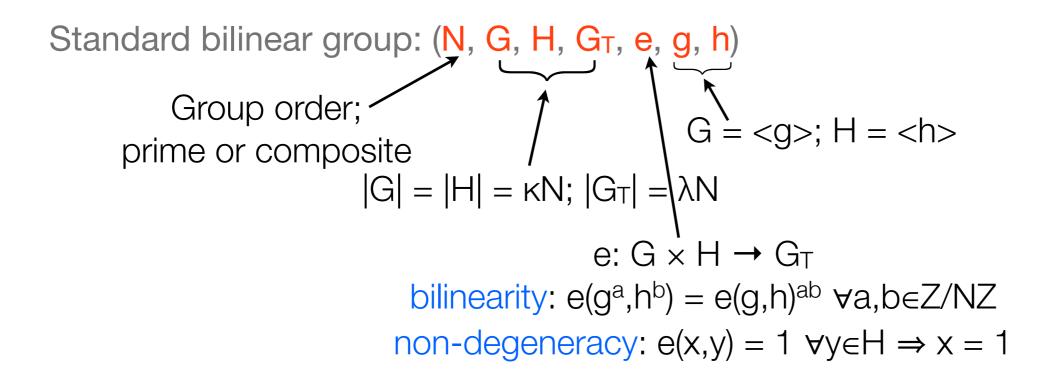
prime or composite

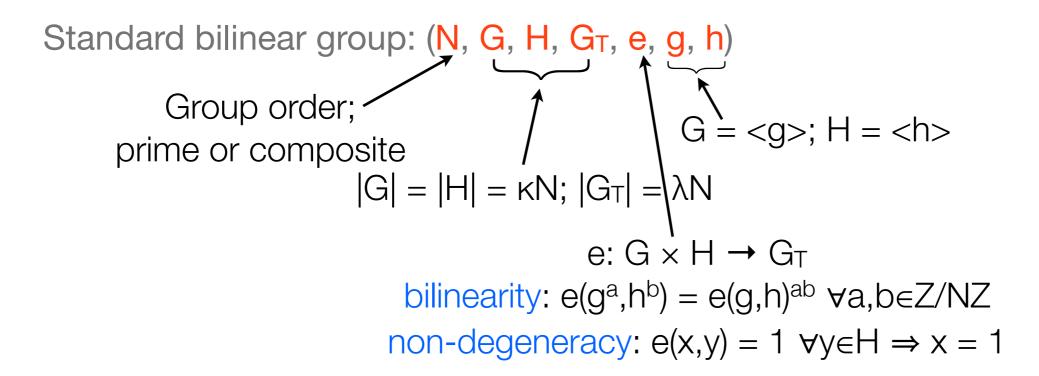
$$|G| = |H| = \kappa N; |G_T| = \lambda N$$

```
Standard bilinear group: (N, G, H, G<sub>T</sub>, e, g, h)

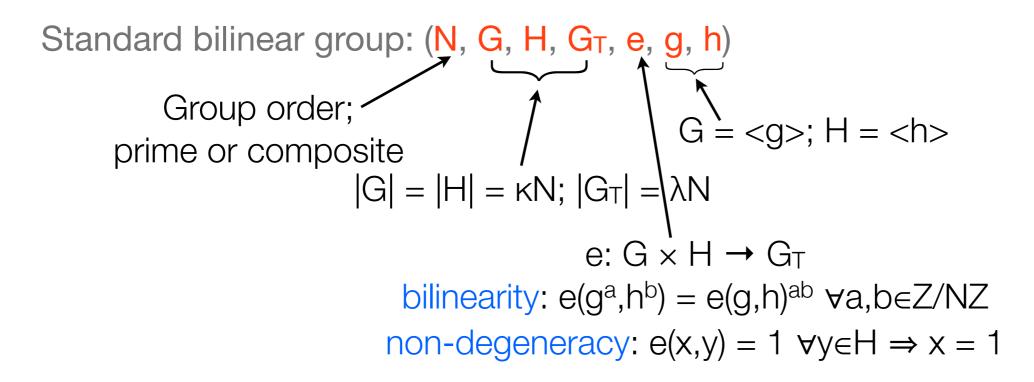
Group order;

prime or composite
|G| = |H| = \kappa N; |G_T| = \lambda N
e: G \times H \to G_T
bilinearity: e(g^a, h^b) = e(g, h)^{ab} \ \forall a, b \in Z/NZ
non-degeneracy: e(x,y) = 1 \ \forall y \in H \Rightarrow x = 1
```

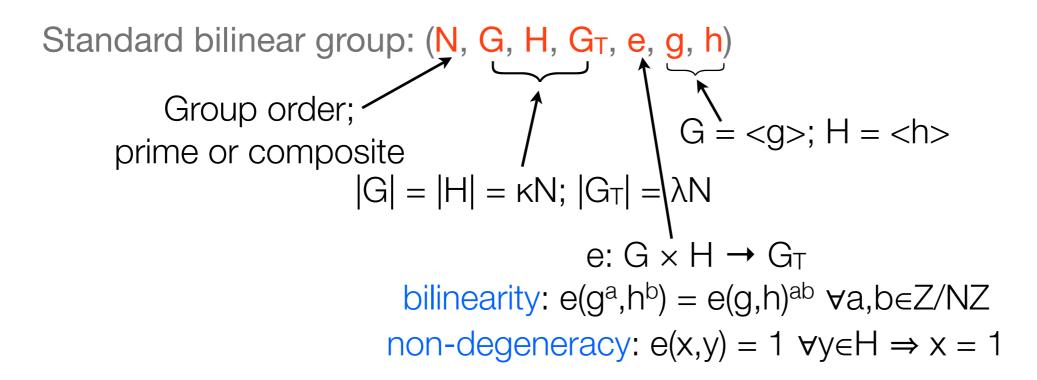






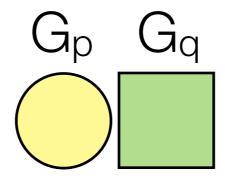


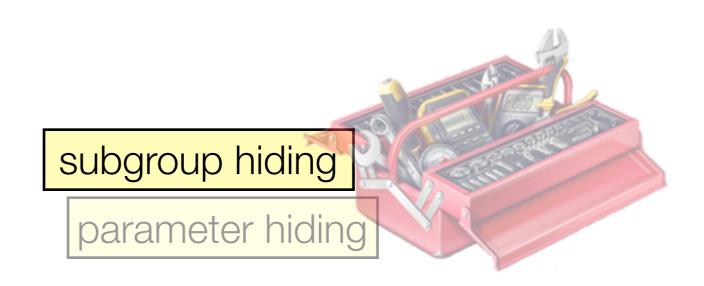




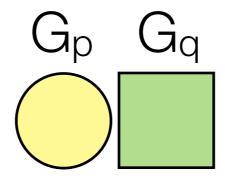








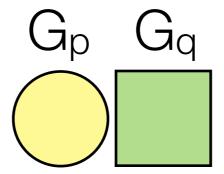
Composite-order bilinear group: (N, G,  $G_T$ , e, g) where N = pq



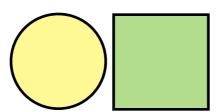
Subgroup hiding [BGN05]:



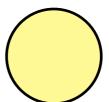
Composite-order bilinear group: (N, G,  $G_T$ , e, g) where N = pq



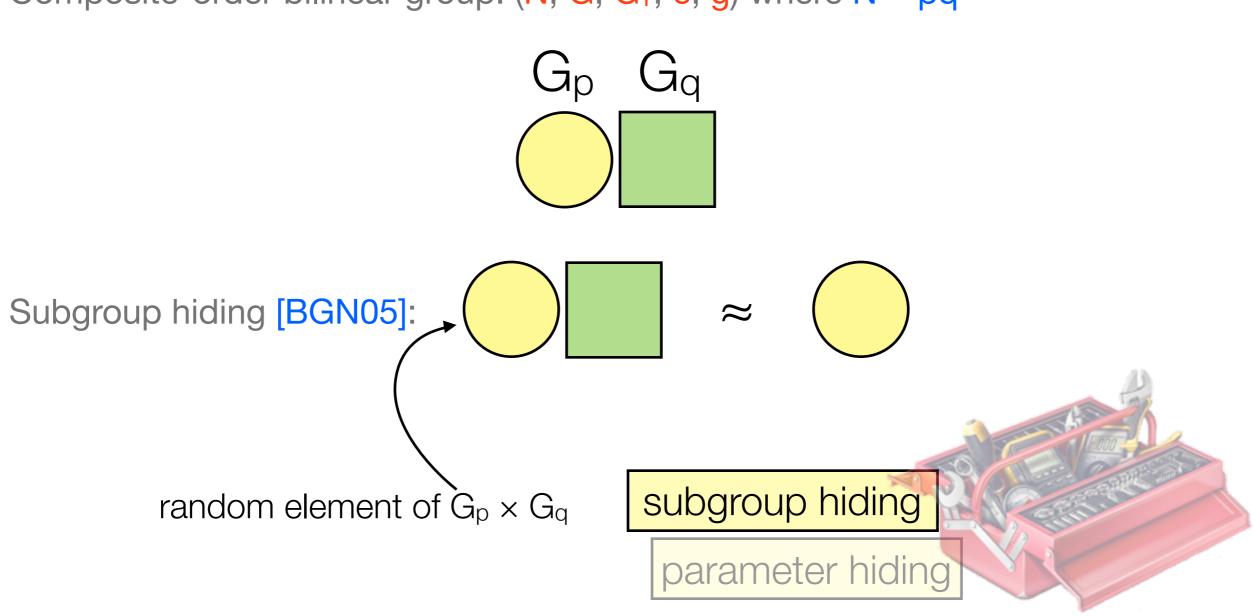
Subgroup hiding [BGN05]:

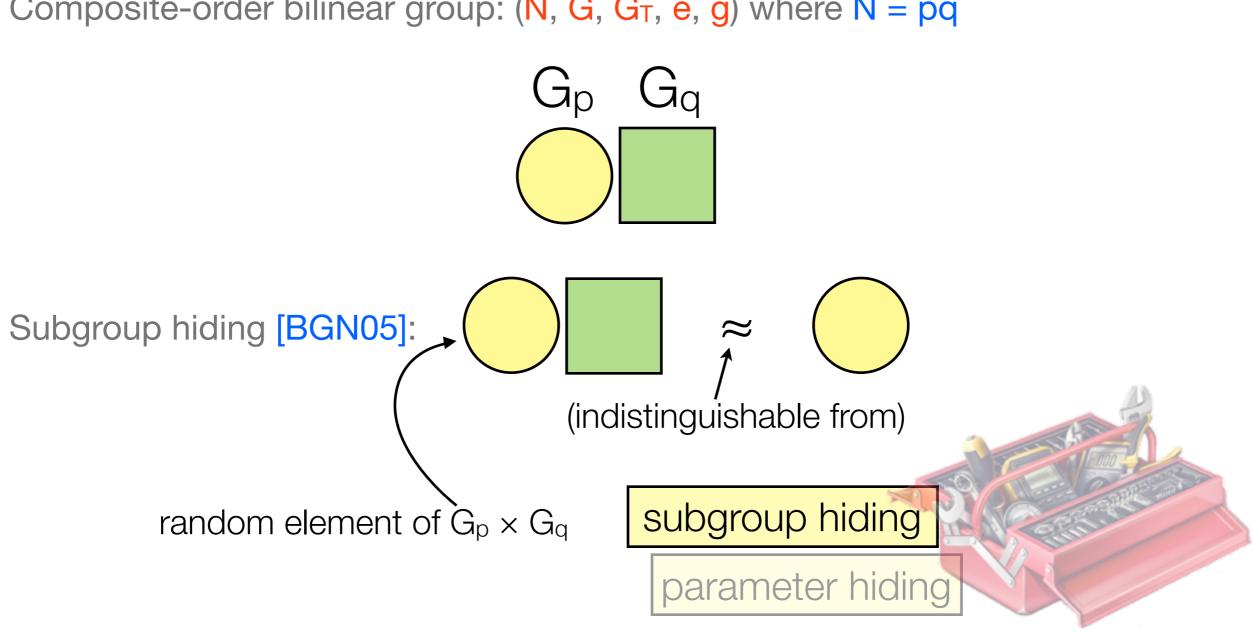


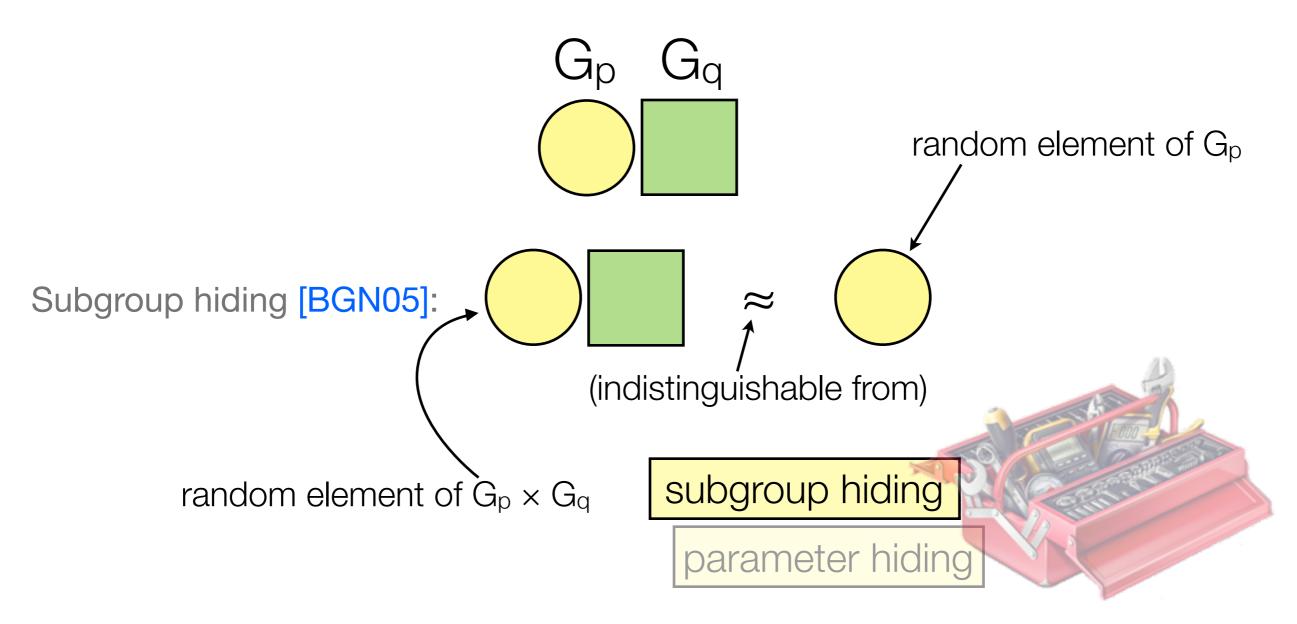
 $\approx$ 



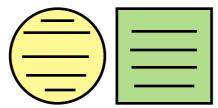
subgroup hiding parameter hiding













$$g_1^{f(x_1,...,x_c)}g_2^{f(x_1,...,x_c)}$$



$$g_1^{f(x_1,\ldots,x_C)}g_2^{f(x_1,\ldots,x_C)}$$



$$g_1^{f(x_1,...,x_c)}g_2^{f(x_1,...,x_c)} \approx g_1^{f(x_1,...,x_c)}g_2^{f(x_1,...,x_c)}$$



Parameter hiding: elements correlated across subgroups are distributed identically to uncorrelated elements

$$g_1^{f(x_1,...,x_C)}g_2^{f(x_1,...,x_C)} \approx g_1^{f(x_1,...,x_C)}g_2^{f(x_1,...,x_C)}$$

is independent from



Parameter hiding: elements correlated across subgroups are distributed identically to uncorrelated elements

$$g_1^{f(x_1,...,x_C)}g_2^{f(x_1,...,x_C)} \approx g_1^{f(x_1,...,x_C)}g_2^{f(x_1',...,x_C')}$$





xi mod p reveals nothing about xi mod q (CRT)

## Challenge ciphertext

### Challenge ciphertext

#### **ID** queries

#### Challenge ciphertext

normal:

**ID** queries

#### Challenge ciphertext

normal:

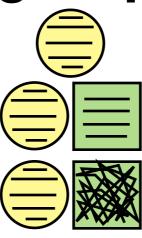


(subgroup hiding)

**ID** queries

#### Challenge ciphertext

normal:



(subgroup hiding)

(parameter hiding)

**ID** queries

# Challenge ciphertext

normal:

semi-functional (SF):

(subgroup hiding)

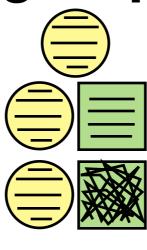
(parameter hiding)

**ID** queries

#### Challenge ciphertext

normal:

semi-functional (SF):



(subgroup hiding)

(parameter hiding)

**ID** queries

normal:

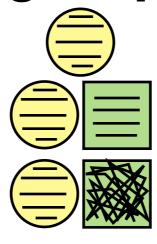


(subgroup hiding)

#### Challenge ciphertext

normal:

semi-functional (SF):



(subgroup hiding)

(parameter hiding)

**ID** queries

normal:



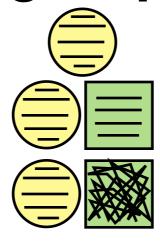
(subgroup hiding)

(parameter hiding)

#### Challenge ciphertext

normal:

semi-functional (SF):



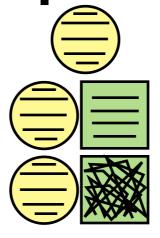
(subgroup hiding)

(parameter hiding)

**ID** queries

normal:

semi-functional (SF):



(subgroup hiding)

(parameter hiding)

#### Challenge ciphertext

normal:

(subgroup hiding)

semi-functional (SF):

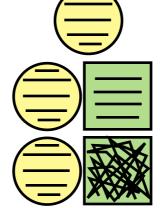


(parameter hiding)

SF keys don't decrypt SF ciphertexts!

#### **ID** queries

normal:



(subgroup hiding)

(parameter hiding)

semi-functional (SF):

1. start with base scheme

normal:



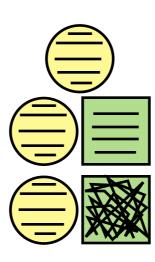
1. start with base scheme



- 1. start with base scheme
- 2. transition to SF version

normal:

semi-functional (SF):



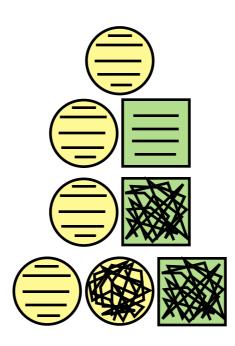
(subgroup hiding)

(parameter hiding)

- 1. start with base scheme
- 2. transition to SF version

normal:

semi-functional (SF):



(subgroup hiding)

(parameter hiding)

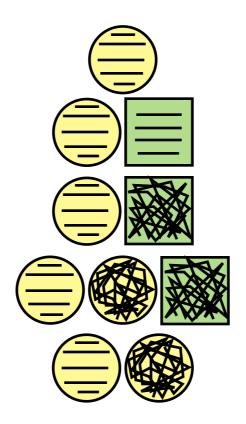
(subgroup hiding)

- 1. start with base scheme
- 2. transition to SF version

#### Dual systems in three easy steps

normal:

semi-functional (SF):



(subgroup hiding)

(parameter hiding)

(subgroup hiding)

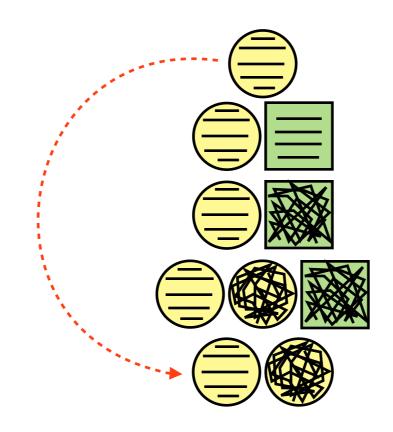
(subgroup hiding)

- 1. start with base scheme
- 2. transition to SF version

#### Dual systems in three easy steps

normal:

semi-functional (SF):



(subgroup hiding)

(parameter hiding)

(subgroup hiding)

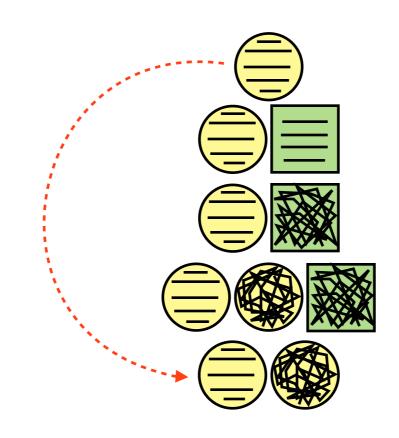
(subgroup hiding)

- 1. start with base scheme
- 2. transition to SF version

#### Dual systems in three easy steps

normal:

semi-functional (SF):



(subgroup hiding)

(parameter hiding)

(subgroup hiding)

(subgroup hiding)

- 1. start with base scheme
- 2. transition to SF version
- 3. argue information is hidden,

#### Outline

Bilinear groups

q-Type assumptions

The uber-assumption
Relating uber-assumptions
A bijection trick

Extensions

Conclusions

Uber-assumption is parameterized by (c,R,S,T,f)

• c = number of variables:  $x_1,...,x_c \leftarrow \mathcal{R}$ 

- c = number of variables:  $x_1,...,x_c \leftarrow \mathcal{R}$
- $R = \langle 1, \rho_1, ..., \rho_r \rangle$ : A is given g,  $\{g^{\rho_i(x_1, ..., x_c)}\}$

- c = number of variables:  $x_1,...,x_c \leftarrow \mathcal{R}$
- $R = \langle 1, \rho_1, ..., \rho_r \rangle$ : A is given g,  $\{g^{\rho_i(x_1, ..., x_c)}\}$
- $S = \langle 1, \sigma_1, ..., \sigma_s \rangle$ : A is given h,  $\{h^{\sigma_i(x_1, ..., x_c)}\}$

- c = number of variables:  $x_1,...,x_c \leftarrow \mathcal{R}$
- $R = \langle 1, \rho_1, ..., \rho_r \rangle$ : A is given g,  $\{g^{\rho_i(x_1, ..., x_c)}\}$
- $S = \langle 1, \sigma_1, ..., \sigma_s \rangle$ : A is given h,  $\{h^{\sigma_i(x_1, ..., x_c)}\}$
- $T = \langle 1, \tau_1, ..., \tau_t \rangle$ : A is given  $e(g,h), \{e(g,h)^{\tau_i(x_1,...,x_c)}\}$

- c = number of variables:  $x_1,...,x_c \leftarrow \mathcal{R}$
- $R = \langle 1, \rho_1, ..., \rho_r \rangle$ : A is given g,  $\{g^{\rho_i(x_1, ..., x_c)}\}$
- $S = \langle 1, \sigma_1, ..., \sigma_s \rangle$ : A is given h,  $\{h^{\sigma_i(x_1, ..., x_c)}\}$
- $T = \langle 1, \tau_1, ..., \tau_t \rangle$ : A is given  $e(g,h), \{e(g,h)^{\tau_i(x_1,...,x_c)}\}$
- $f(x_1,...,x_c)$ : A needs to compute  $e(g,h)^{f(x_1,...,x_c)}$  (or distinguish it from random)

Uber-assumption is parameterized by (c,R,S,T,f)

- c = number of variables:  $x_1,...,x_c \leftarrow \mathcal{R}$
- $R = \langle 1, \rho_1, ..., \rho_r \rangle$ : A is given g,  $\{g^{\rho_i(x_1, ..., x_c)}\}$
- $S = \langle 1, \sigma_1, ..., \sigma_s \rangle$ : A is given h,  $\{h^{\sigma_i(x_1, ..., x_c)}\}$
- $T = \langle 1, \tau_1, ..., \tau_t \rangle$ : A is given  $e(g,h), \{e(g,h)^{\tau_i(x_1,...,x_c)}\}$
- $f(x_1,...,x_c)$ : A needs to compute  $e(g,h)^{f(x_1,...,x_c)}$  (or distinguish it from random)

uber(c,R,S,T,f) assumption: given (R,S,T) values, hard to compute/distinguish f

exponent q-SDH [ZS-NS04]: given (g,gx,...,gxq), distinguish gxq+1 from random

• c = number of variables: c = 1

- c = number of variables: c = 1
- $R = \langle 1, \rho_1, ..., \rho_r \rangle$ :  $\rho_i(x) = x^i \ (\forall i \ 0 \le i \le q)$

- c = number of variables: c = 1
- $R = \langle 1, \rho_1, ..., \rho_r \rangle$ :  $\rho_i(x) = x^i \ (\forall i \ 0 \le i \le q)$
- S = <1>
- T = <1>

- c = number of variables: c = 1
- $R = \langle 1, \rho_1, ..., \rho_r \rangle$ :  $\rho_i(x) = x^i \ (\forall i \ 0 \le i \le q)$
- S = <1>
- T = <1>
- $f(x_1,...,x_c)$ :  $f(x) = x^{q+1}$

exponent q-SDH [ZS-NS04]: given (g,gx,...,gxq), distinguish gxq+1 from random

- c = number of variables: c = 1
- $R = \langle 1, \rho_1, ..., \rho_r \rangle$ :  $\rho_i(x) = x^i \ (\forall i \ 0 \le i \le q)$
- S = <1>
- T = <1>
- $f(x_1,...,x_c)$ :  $f(x) = x^{q+1}$

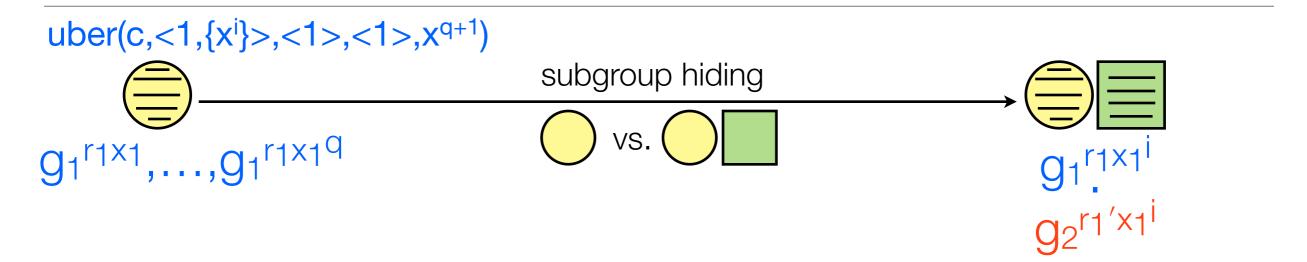
exponent q-SDH is uber(1,<1, $\{x^i\}$ >,<1>,<1>, $x^{q+1}$ )

uber(c,<1, $\{x^i\}$ >,<1>,<1>, $x^{q+1}$ )

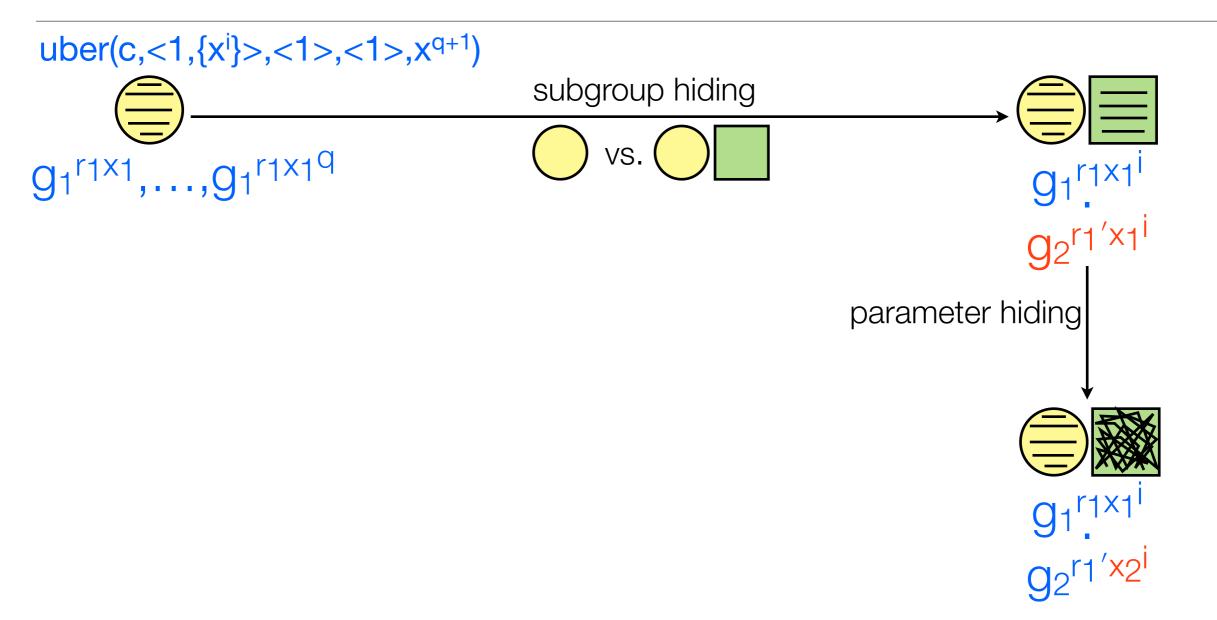
- 1. start with base scheme
- 2. transition to SF version
- 3. argue information is hidden

- 1. start with base scheme
- 2. transition to SF version
- 3. argue information is hidden

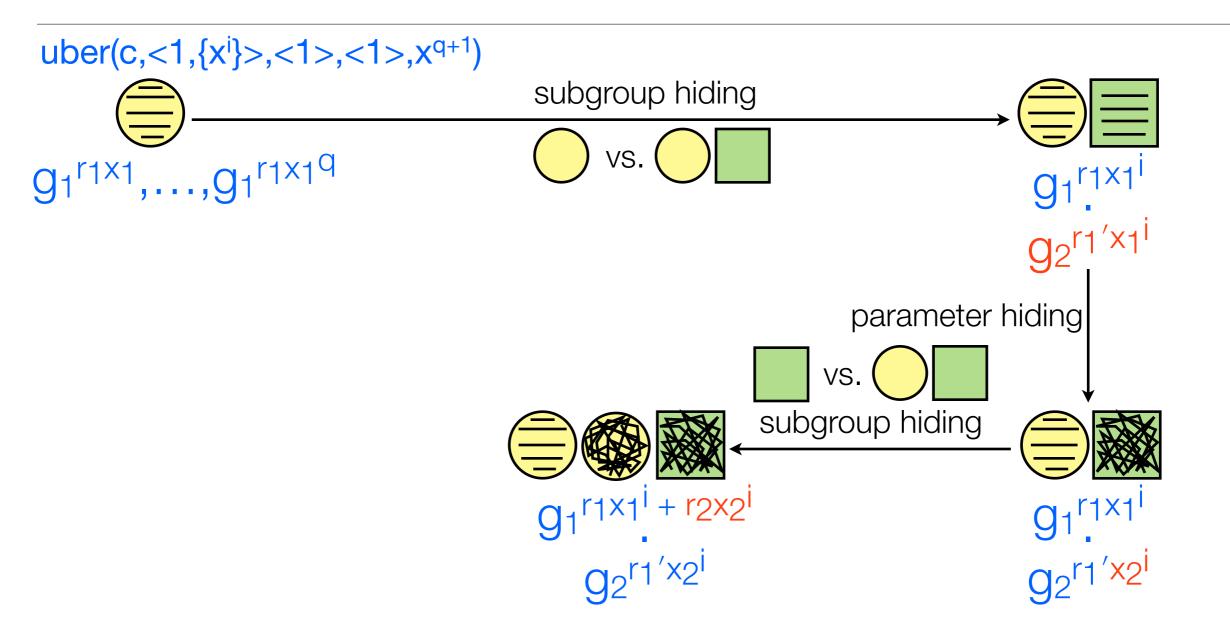
- 1. start with base scheme
- 2. transition to SF version
- 3. argue information is hidden



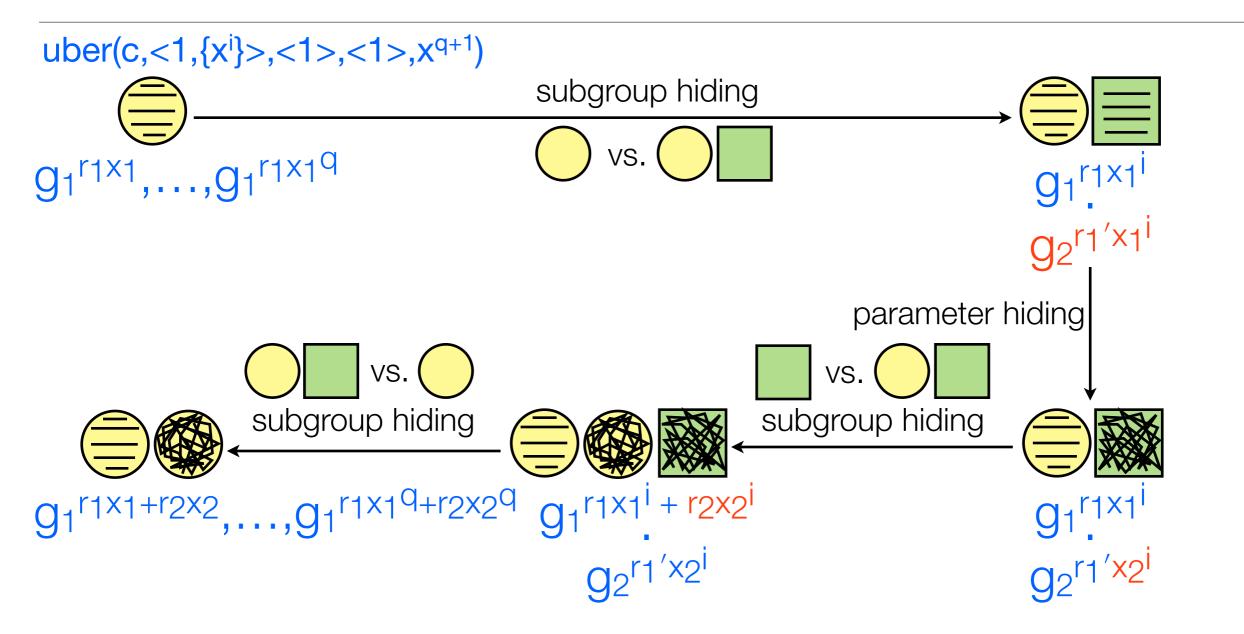
- 1. start with base scheme
- 2. transition to SF version
- 3. argue information is hidden



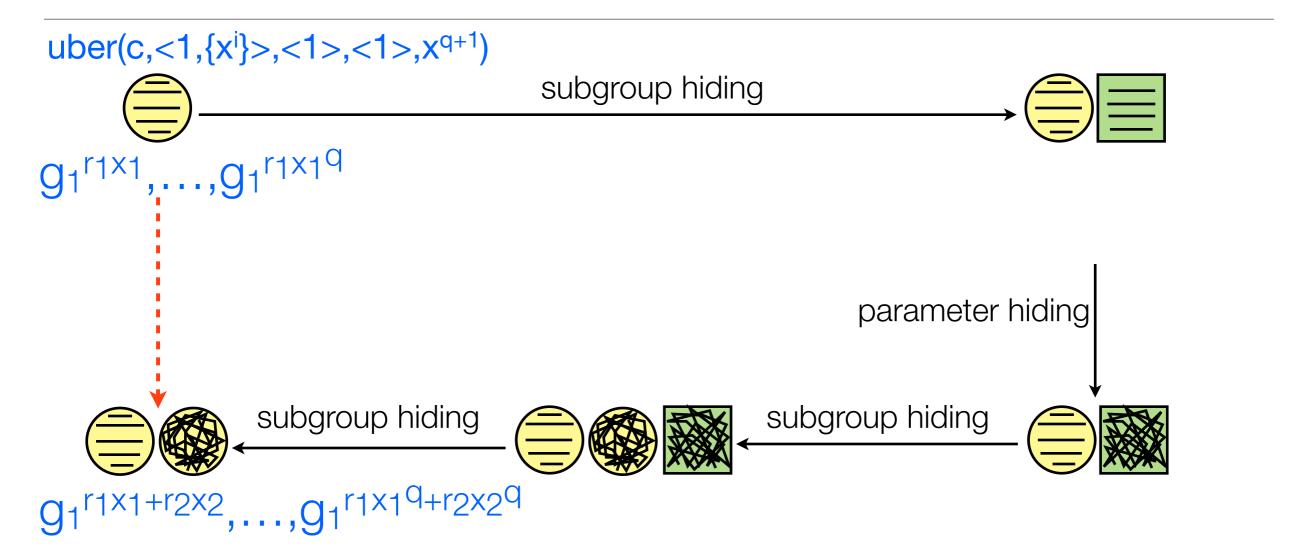
- 1. start with base scheme
- 2. transition to SF version
- 3. argue information is hidden



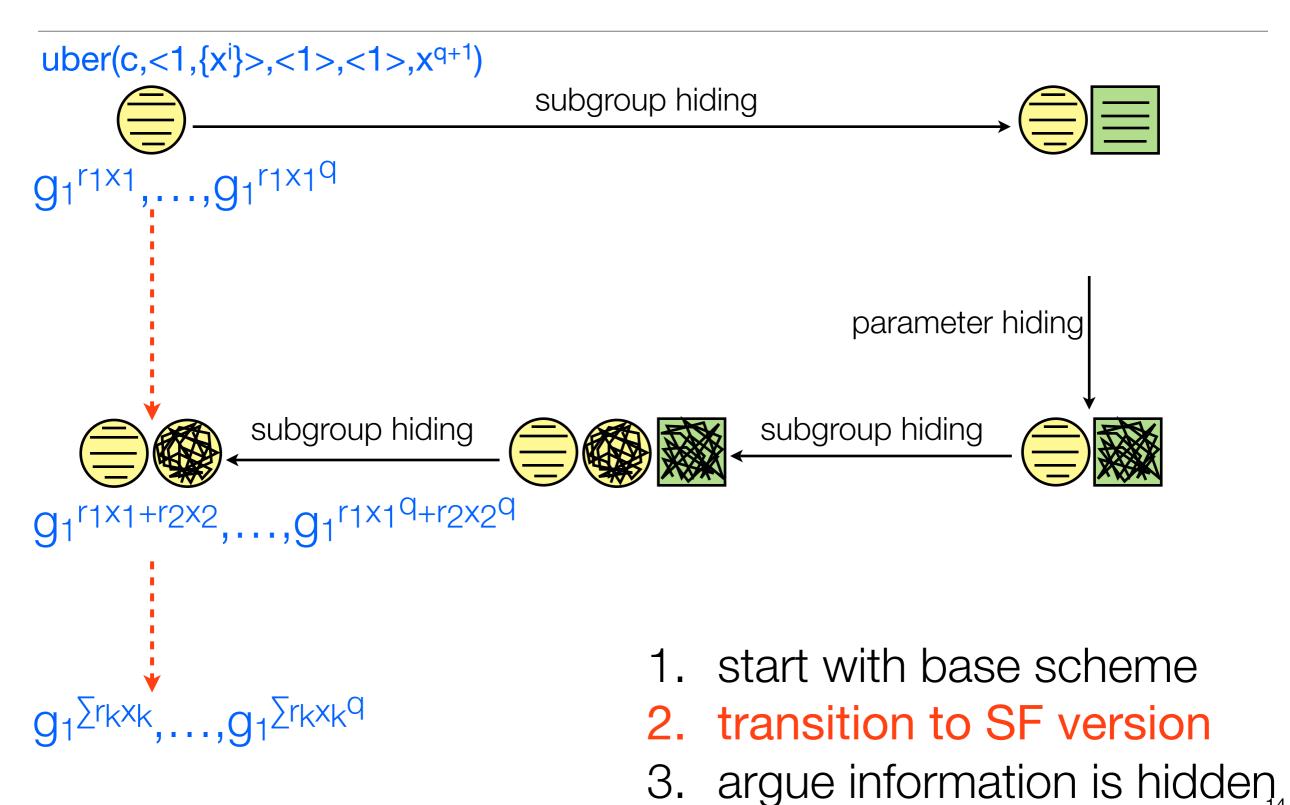
- 1. start with base scheme
- 2. transition to SF version
- 3. argue information is hidden

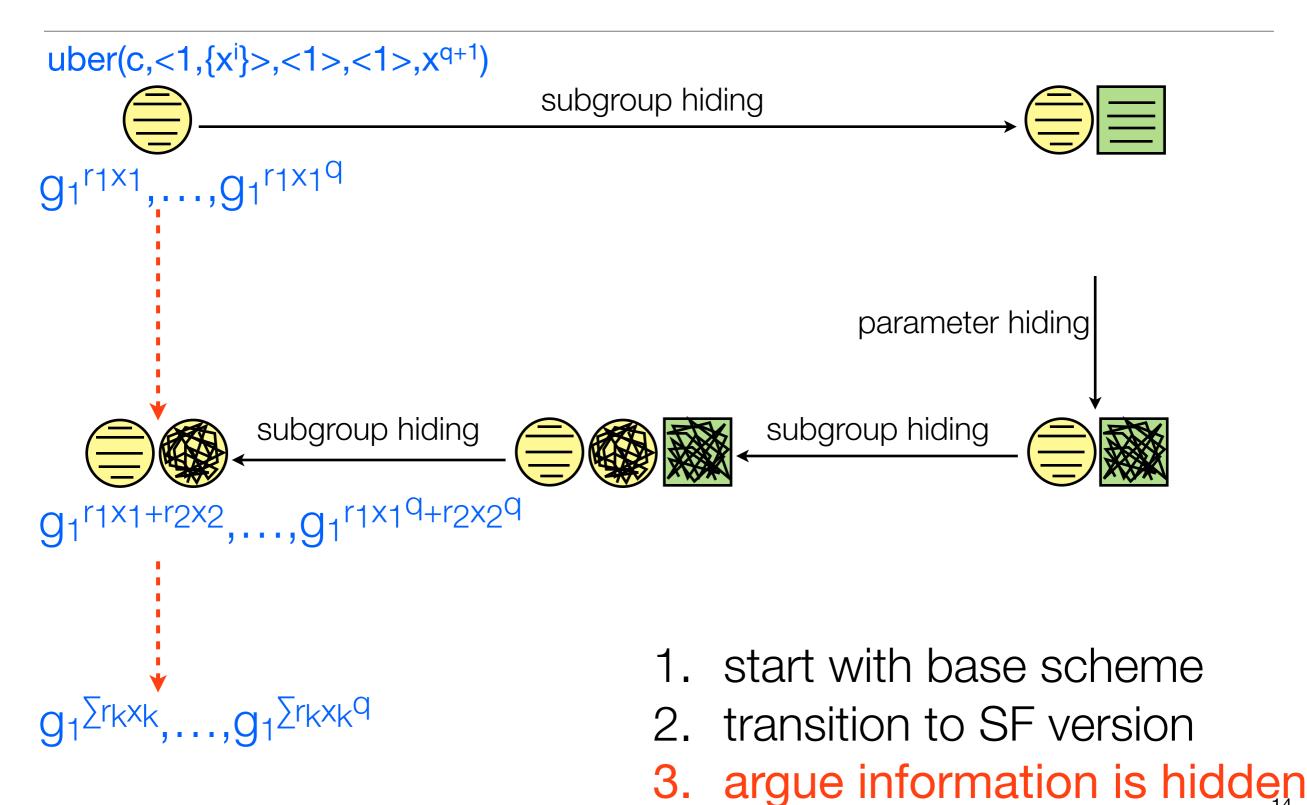


- 1. start with base scheme
- 2. transition to SF version
- 3. argue information is hidden



- 1. start with base scheme
- transition to SF version
- argue information is hidden,





- 1. start with base scheme
- 2. transition to SF version
- 3. argue information is hidden

uber(c,R,<1, $\{x^i\}$ >,<1>, $x^{q+1}$ )  $\rightarrow$  uber( $\ell$ c,<1, $\{\sum r_k x_k^i\}$ >,<1>,<1>, $\sum r_k x_k^{q+1}$ )

- 1. start with base scheme
- 2. transition to SF version
- 3. argue information is hidden

uber(c,R,<1, $\{x^i\}$ >,<1>, $x^{q+1}$ )  $\rightarrow$  uber( $\ell$ c,<1, $\{\sum r_k x_k^i\}$ >,<1>,<1>, $\sum r_k x_k^{q+1}$ )

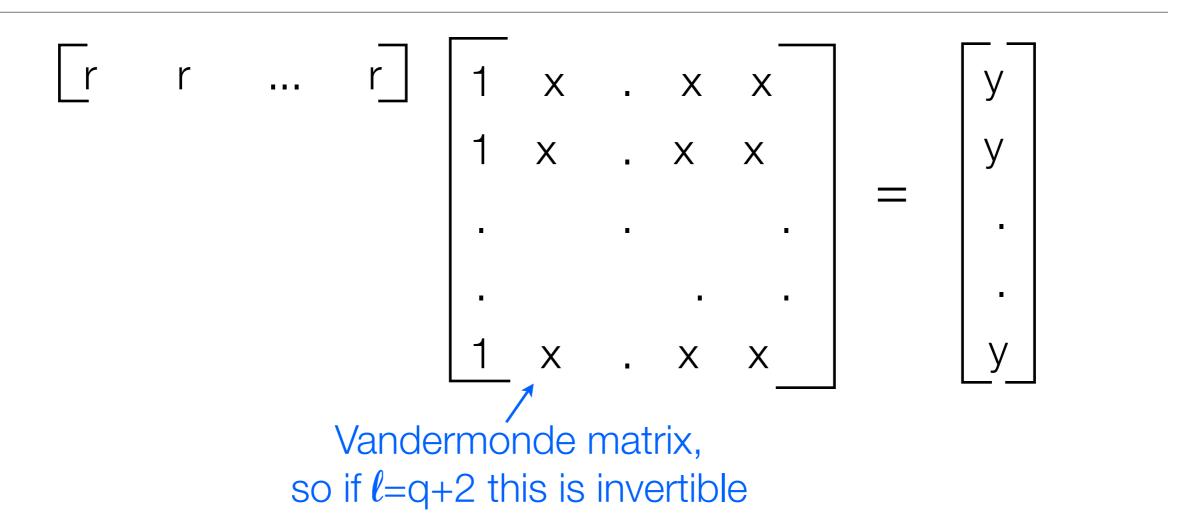
- 1. start with base scheme
- 2. transition to SF version
- 3. argue information is hidden

 $uber(c,R,<1,\{x^i\}>,<1>,x^{q+1}) \rightarrow uber(\ell c,<1,\{\sum r_k x_k{}^i\}>,<1>,<1>,\sum r_k x_k{}^{q+1})$ 

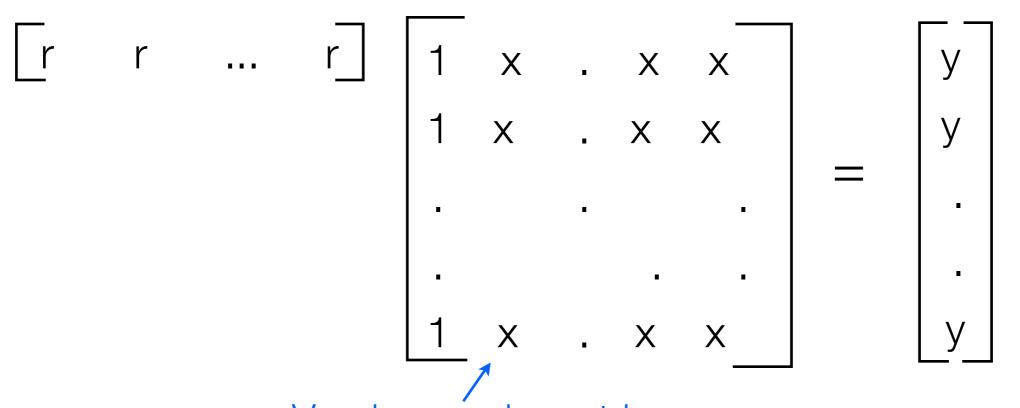
So A is really given

- 1. start with base scheme
- 2. transition to SF version
- 3. argue information is hidden

- start with base scheme
- 2. transition to SF version
- 3. argue information is hidden



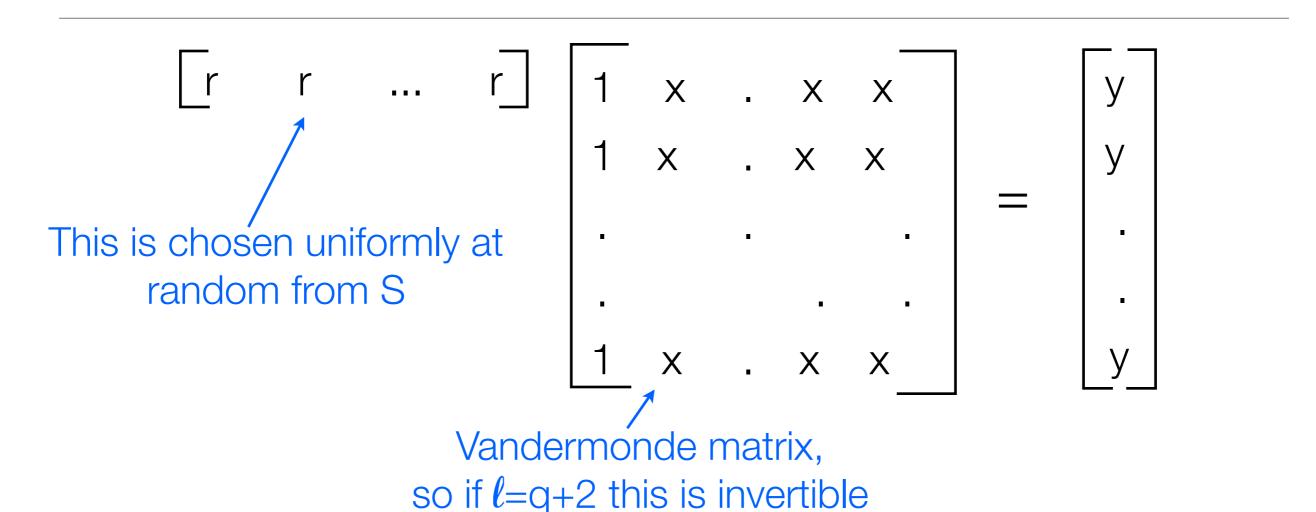
- 1. start with base scheme
- 2. transition to SF version
- 3. argue information is hidden



Vandermonde matrix, so if  $\ell$ =q+2 this is invertible

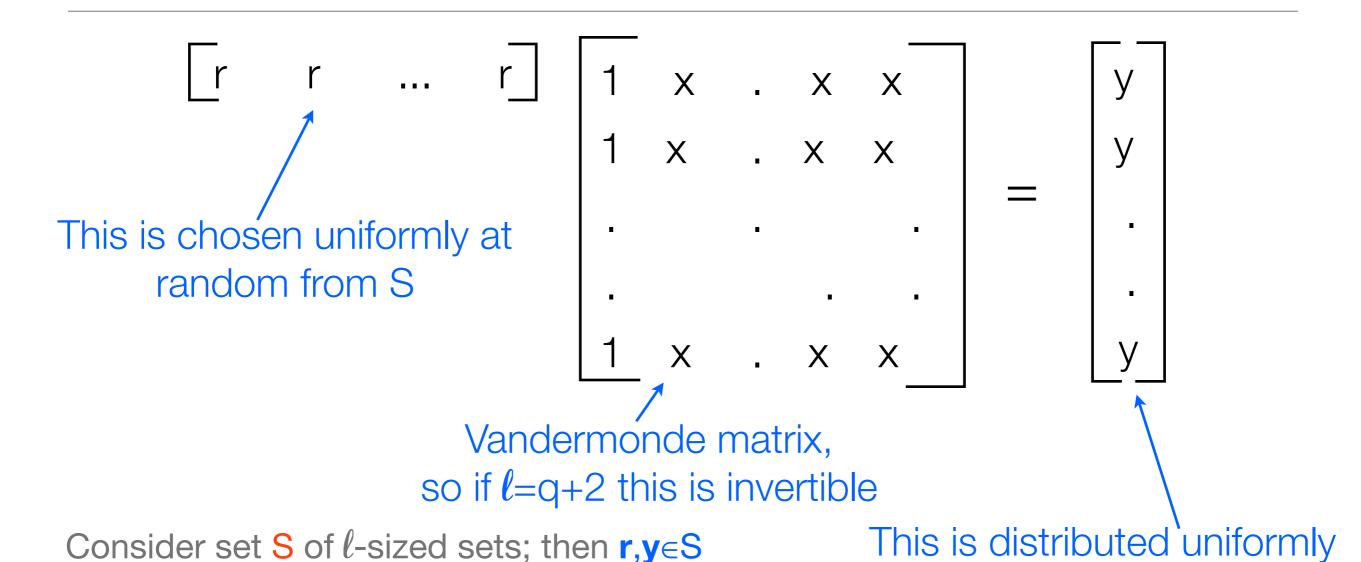
Consider set S of  $\ell$ -sized sets; then  $r,y \in S$ 

- permutation
  Matrix multiplication is <del>map</del> M: S → S
- 1. start with base scheme
- 2. transition to SF version
- 3. argue information is hidden



Consider set S of  $\ell$ -sized sets; then  $r,y \in S$ 

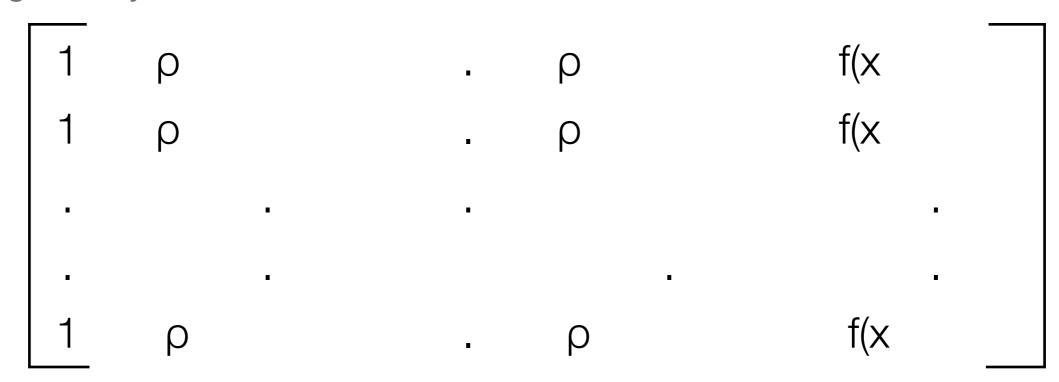
- permutation
  Matrix multiplication is <del>map</del> M: S → S
- 1. start with base scheme
- 2. transition to SF version
- 3. argue information is hidden



- permutation
- Matrix multiplication is <del>map</del> M: S → S
- 1. start with base scheme
  - 2. transition to SF version
  - 3. argue information is hidden

random as well!

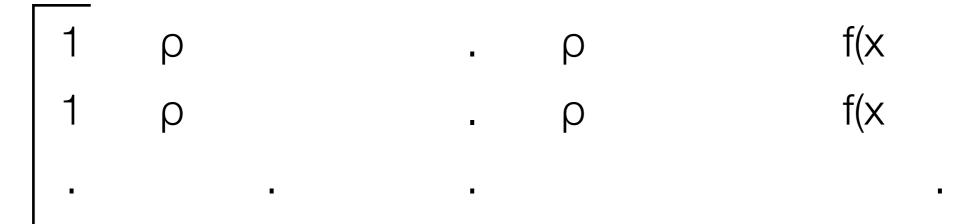
More generally, this is true if



has linearly independent columns (or rows)

- 1. start with base scheme
- 2. transition to SF version
- 3. argue information is hidden

More generally, this is true if



has lir

Decisional uber(c,R,S,T,f) holds if:

1. subgroup hiding and parameter hiding hold 2. S = T = <1>

3. f is not a linear combination of pi

More generally, this is true if

 $1~\rho$  .  $\rho$  f(x

1 ρ . ρ f(x

- only computational requirement

has lir

Decisional uber(c,R,S,T,f) holds if:

1. subgroup hiding and parameter hiding hold

2. 
$$S = T = \langle 1 \rangle$$

3. f is not a linear combination of pi

More generally, this is true if

1 ρ . ρ f(x 1 ρ . ρ

- only computational requirement

has lir

Decisional uber(c,R,S,T,f) holds if:

1. subgroup hiding and parameter hiding hold

limitation 
$$\longrightarrow$$
 2. S = T = <1>

3. f is not a linear combination of pi

<del>o. argue imormation is n</del>iddeņ

#### Outline

Bilinear groups

q-Type assumptions

#### Extensions

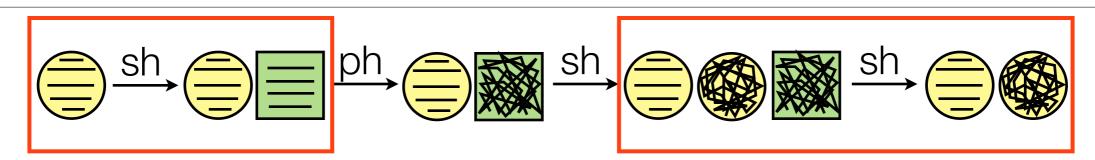
Broader classes of assumptions Dodis-Yampolskiy PRF Conclusions





Remember that we needed two types of subgroup hiding ...

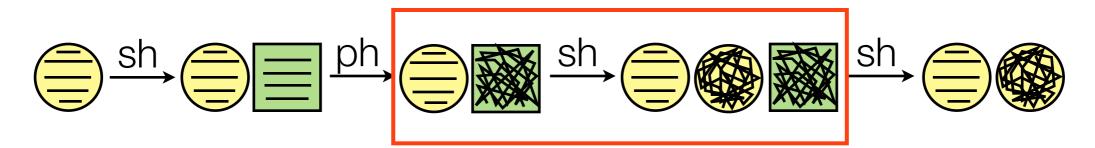
...even when given a generator for



Remember that we needed two types of subgroup hiding ...



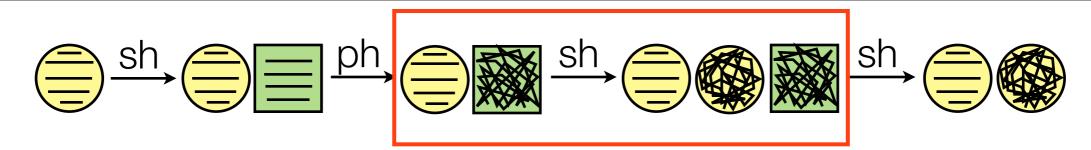
...even when given a generator for



Remember that we needed two types of subgroup hiding ...



...even when given a generator for



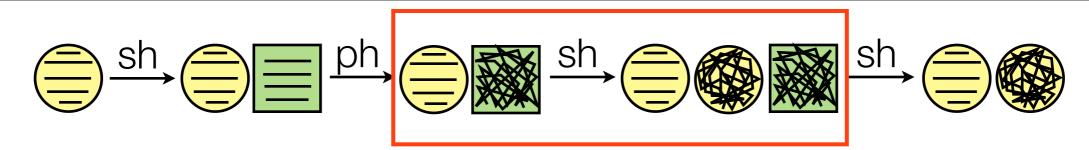
Remember that we needed two types of subgroup hiding ...



...even when given a generator for

This restricts us to "one-sided" assumptions 2.S = T = <1>

2. 
$$S = T = \langle 1 \rangle$$



Remember that we needed two types of subgroup hiding ...

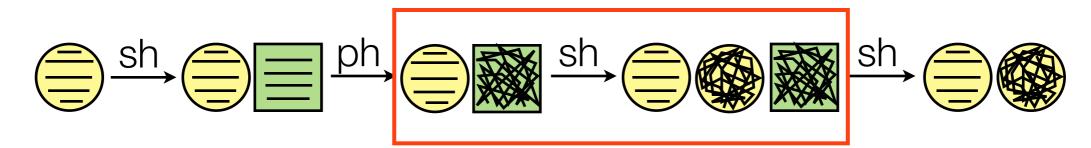


...even when given a generator for

This restricts us to "one-sided" assumptions 2.S = T = <1>

2. 
$$S = T = <1>$$

$$(g,g^x,...,g^{x^q}) \rightarrow g^{x^{q+1}}$$
 or random



Remember that we needed two types of subgroup hiding ...





...even when given a generator for

This restricts us to "one-sided" assumptions 2.S = T = <1>

2. 
$$S = T = <1>$$

$$(g,g^x,...,g^{x^q}) \rightarrow g^{x^{q+1}}$$
 or random

$$(g,g^x,...,g^{x^q},h^x) \rightarrow compute (c,g^{1/x+c})$$



Remember that we needed two types of subgroup hiding...







Remember that we needed two types of subgroup hiding...



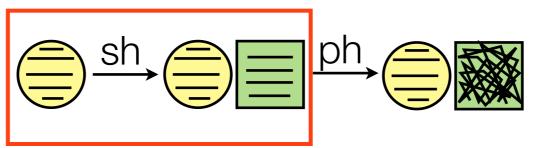
To address this, switch back to regular dual systems



Remember that we needed two types of subgroup hiding...



To address this, switch back to regular dual systems

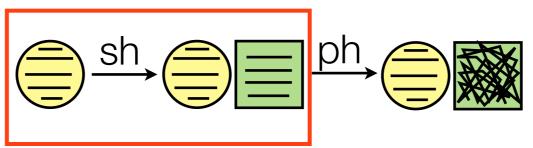




Remember that we needed two types of subgroup hiding...



To address this, switch back to regular dual systems

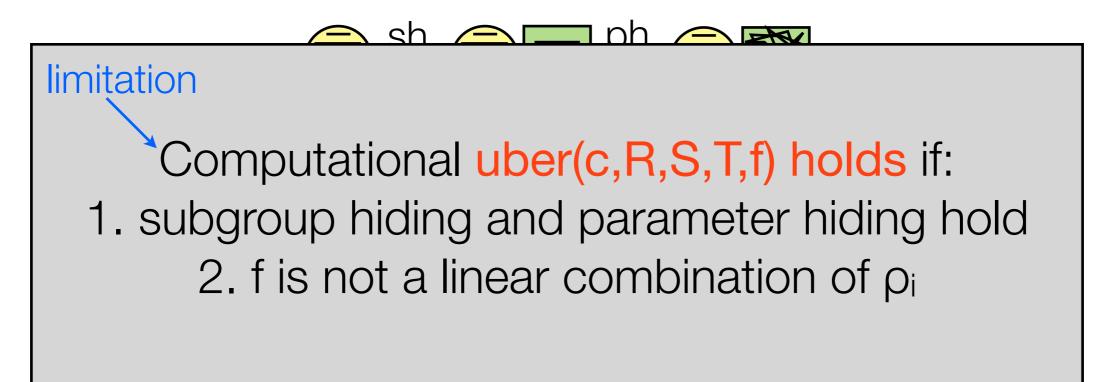




Remember that we needed two types of subgroup hiding...









Remember that we needed two types of subgroup hiding...



To address this, switch back to regular dual systems

This implies (for example) that q-SDH [BB04] follows from subgroup hiding....

...and so does everything based on q-SDH (like Boneh-Boyen signatures)\*

\*when instantiated in asymmetric composite-order groups [BRS11]

$$f(x) = u^{1/sk+x}$$
 for fixed  $sk \leftarrow \Re$ ;  $x \in a(\lambda)$ 

 $f(x) = u^{1/sk+x}$  for fixed  $sk \leftarrow \Re$ ;  $x \in a(\lambda)$ 

Theorem [DY05]:  $Adv^{vrf} \le a(\lambda) \cdot Adv^{a(\lambda)-DBDHI}$ 

 $f(x) = u^{1/sk+x}$  for fixed  $sk \leftarrow \Re$ ;  $x \in a(\lambda)$ 

Theorem [DY05]:  $Adv^{vrf} \leq a(\lambda) \cdot Adv^{a(\lambda)-DBDHI}$ 

everifiable random function

$$f(x) = u^{1/sk+x}$$
 for fixed  $sk \leftarrow \Re$ ;  $x \in a(\lambda)$ 

Theorem [DY05]:  $Adv^{vrf} \leq a(\lambda) \cdot Adv^{a(\lambda)-DBDHI}$ 

- verifiable random function
- erequire u=e(g,h)

$$f(x) = u^{1/sk+x}$$
 for fixed  $sk \leftarrow \Re$ ;  $x \in a(\lambda)$ 

Theorem [DY05]: 
$$Adv^{vrf} \leq a(\lambda) \cdot Adv^{a(\lambda)-DBDHI}$$

- verifiable random function query assumption
- erequire u=e(g,h)

$$f(x) = u^{\frac{1}{sk+x}}$$
 for fixed  $sk \leftarrow \Re$ ;  $x \in a(\lambda)$ 

Theorem [DY05]: 
$$Adv^{vrf} \leq a(\lambda) - Adv^{a(\lambda)-DBDHI}$$

- verifiable random function q-type assumption
- $\cong$  require u=e(g,h) = looseness: need  $|a(\lambda)| \le poly(\lambda)$

$$f(x) = u^{1/sk+x}$$
 for fixed  $sk \leftarrow \Re$ ;  $x \in a(\lambda)$ 

Theorem [DY05]:  $Adv^{vrf} \le a(\lambda) \cdot Adv^{a(\lambda)-DBDHI}$ 

verifiable random function query assumption

 $\stackrel{\smile}{=}$  require u=e(g,h)  $\stackrel{\smile}{=}$  looseness: need  $|a(\lambda)| \le poly(\lambda)$ 

Theorem: Adv<sup>prf</sup> ≤ q · Adv<sup>sgh</sup>

$$f(x) = u^{1/sk+x}$$
 for fixed  $sk \leftarrow \Re$ ;  $x \in a(\lambda)$ 

Theorem [DY05]:  $Adv^{vrf} \le a(\lambda) \cdot Adv^{a(\lambda)-DBDHI}$ 

- verifiable random function query assumption
- $\cong$  require u=e(g,h) = looseness: need  $|a(\lambda)| \le poly(\lambda)$

Theorem: Adv<sup>prf</sup> ≤ q · Adv<sup>sgh</sup>

pseudorandom function

$$f(x) = u^{1/sk+x}$$
 for fixed  $sk \leftarrow \Re$ ;  $x \in a(\lambda)$ 

Theorem [DY05]:  $Adv^{vrf} \le a(\lambda) \cdot Adv^{a(\lambda)-DBDHI}$ 

- verifiable random function q-type assumption

require u=e(g,h)

 $|a(\lambda)| \leq poly(\lambda)$ 

Theorem: Adv<sup>prf</sup> ≤ q · Adv<sup>sgh</sup>

- pseudorandom function
- "require composite order

$$f(x) = u^{1/sk+x}$$
 for fixed  $sk \leftarrow \Re$ ;  $x \in a(\lambda)$ 

Theorem [DY05]:  $Adv^{vrf} \le a(\lambda) \cdot Adv^{a(\lambda)-DBDHI}$ 

verifiable random function q-type assumption

require u=e(g,h)

 $|a(\lambda)| \leq poly(\lambda)$ 

Theorem: Advoprf ≤ q · Advosgh

pseudorandom function
static assumption

"require composite order

$$f(x) = u^{1/sk+x}$$
 for fixed  $sk \leftarrow \Re$ ;  $x \in a(\lambda)$ 

Theorem [DY05]:  $Adv^{vrf} \leq a(\lambda) \cdot Adv^{a(\lambda)-DBDHI}$ 

- verifiable random function q-type assumption

require u=e(g,h)

 $|a(\lambda)| \leq poly(\lambda)$ 

Theorem: Adv<sup>prf</sup> ≤ q Adv<sup>sgh</sup>

- pseudorandom function
  static assumption
- "require composite order
- (a) of arbitrary size

#### Outline

Bilinear groups q-Type assumptions

Extensions

Conclusions

We applied the dual-system technique directly to a broad class of assumptions

We applied the dual-system technique directly to a broad class of assumptions

Limitation: Restricted to (asymmetric) composite-order (bilinear) groups

We applied the dual-system technique directly to a broad class of assumptions

Limitation: Restricted to (asymmetric) composite-order (bilinear) groups

Limitation: Can't get rid of every q-type assumption

We applied the dual-system technique directly to a broad class of assumptions

Limitation: Restricted to (asymmetric) composite-order (bilinear) groups

Limitation: Can't get rid of every q-type assumption

Full version!: cs.ucsd.edu/~smeiklejohn/files/eurocrypt14a.pdf

We applied the dual-system technique directly to a broad class of assumptions

Limitation: Restricted to (asymmetric) composite-order (bilinear) groups

Limitation: Can't get rid of every q-type assumption

Full version!: cs.ucsd.edu/~smeiklejohn/files/eurocrypt14a.pdf

Thanks! Any questions?