Flexible Models for Secure Systems

Sarah Meiklejohn











Motivating scenario that founded modern cryptography:



Security model for this interaction is well established

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Encryption works for secure online communication: SSL/TLS [1996]

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Problem #1: everyone is an adversary!



3

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3

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Problem #2: what are the security goals?



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Problem #2: what are the security goals?



can Bob retrieve Alice's files? can Dropbox read Alice's files?





Problem #2: what are the security goals?



- can Bob retrieve Alice's files?
- can Dropbox read Alice's files?
- did Dropbox delete Alice's files?





Problem #2: what are the security goals?



can Bob retrieve Alice's files?
can Dropbox read Alice's files?
did Dropbox delete Alice's files?
can Dropbox efficiently store files?





Problem #2: what are the security goals?



can Bob retrieve Alice's files?
can Dropbox read Alice's files?
did Dropbox delete Alice's files?
can Dropbox efficiently store files?
?















theory





outsourced comp. [GGP10]



theory



RKA [B93,K93] MPC [Y82] FHE [G09] outsourced comp. [GGP10] e-cash [C82] ORAM [G096]

theory

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Wickr

theory

practice



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Wickr

theory





RKA [B93,K93] MPC [Y82] FHE [G09] outsourced comp. [GGP10] e-cash [C82] ORAM [G096]

theory

practice

Wickr

diaspora*



RKA [B93,K93]	Bitcoin	SSL/TLS
FHE [G09]	TrueCrypt	Wickr
e-cash [C82] ORAM [GO96]	Tor diaspora*	PGP
theory		practice

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The server's host key is not cached in the registry. You have no guarantee that the server is the computer you think it is. The server's rsa2 key fingerprint is:

ssh-rsa 1024 0a:27:d5:4f:00:9c:d1:a3:ff:ad:5c:cd:b3:7c:83:42

If you trust this host, hit Yes to add the key to

PuTTY's cache and carry on connecting.

If you want to carry on connecting just once, without adding the key to the cache, hit No.

If you do not trust this host, hit Cancel to abandon the connection.

17ijmE4nZstNnevbHLawnhCW9nZ 0.37 BTC 1cTHK2B4D3qg4kj44wxfeQCTE6N 3.11 BTC 1ZhKkzRLLfpNGNHf8qDDH7uRZta 0.20 BTC 1QJ3FK8grGXKEPqVUC1tx6G5Hw 5.16 BTC



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-----BEGIN PGP SIGNATURE-----Version: GnuPG v1.4.11 (GNU/Linux) Comment: Using GnuPG with Thunderbird - http://www.enigmail.net/

iQEcBAEBAgAGBQJSbpK5AAoJEKAl9HCTr8nwh8sIAIIViE5OmpHbzpcatKMdUL9q 0ehAanFyH9WEVxHuv/DAW9Mq35hsYRJL+2wDBUbyeCtClhjgfW0jmgsYc0MevQva w0MUOIeYwBOVkDM6l6E5KN+fRC5qqjeLZvFO2yVrEdItCZeRvNzKm1MREbKnSBPd vkMAW/G/U2dA9GmQjm3lw6Av8Sakh8FC+nGGOwtKHI3TpWK41RdqR503M7tBD1Ov +d9zRchtD5Rs2X43CSTXq5g02zIBcCqarfLKGvPLHOc+4YW+9+9H9V/LT7vTDIDe WWvIrYWM21p8eShdcoukW3aar3ovNqQ6B/8xzLebIAcruvII3ejE2xAR20Btm50= =jT/5

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Signature Algorithm Parameters Not Valid Before Not Valid After Public Key Info Algorithm Parameters	SHA-1 with RSA Encryption (1 2 840 113549 1 1 5) none Wednesday, July 20, 2011 5:00:00 PM PT Thursday, July 18, 2013 4:59:59 PM PT RSA Encryption (1 2 840 113549 1 1 1) none		BEGIN PGP SIGNATURE Version: GnuPG v1.4.11 (GNU/Linux) Comment: Using GnuPG with Thunderbird - http://www.enigmail.net/ iQEcBAEBAgAGBQJSbpK5AAoJEKAI9HCTr8nwh8sIAIIViE5OmpHbzpcatKMdUL9q 0ehAanFyH9WEVxHuv/DAW9Mq35hsYRJL+2wDBUbyeCtClhjgfW0jmgsYc0MevQva w0MUOIeYwBOVkDM6l6E5KN+fRC5qqjeLZvFO2yVrEdItCZeRvNzKm1MREbKnSBPd vkMAW/G/U2dA9GmQjm3lw6Av8Sakh8FC+nGGOwtKHI3TpWK41RdqR503M7tBD10v +d9zRchtD5Rs2X43CSTXq5g02zIBcCqarfLKGvPLHOc+4YW+9+9H9V/LT7vTDIDe WWvIrYWM21p8eShdcoukW3aar3ovNqQ6B/8xzLebIAcruvII3ejE2xAR20Btm50= =jT/5 END PGP SIGNATURE	
Public Key Exponent	8 bytes : D5 14 57 A0 96 40 9F 84 08 C6 66 8D EC E3 03 B2 66 85 AC 5D BB 1C EF 15 93 FD 1F A7 10 49 24 5B 39 D2 60 C8 9A DC C0 CE 40 E6 59 95 B6 52 50 FE 08 25 45 57 73 5F 3A ED DF B6 5C 9D 8A 9C 62 FD 0F 61 BE DC F6 87 80 9D 7F 7C 17 13 77 64 3C 47 F3 87 24 1F 61 E0 81 11 46 E4 DC 50 5C 39 53 E6 68 3D 3C 55 87 C8 BE FC 87 13 D9 5A AA 5D CC 3F C1 74 CD C2 5E 27 16 11	17ijmE4nZstNnevbHLawnhCW9nZ 0.37 BTC 1cTHK2B4D3qg4kj44wxfeQCTE6N 3.11 BTC 1ZhKkzRLLfpNGNHf8qDDH7uRZta 0.20 BTC 1QJ3FK8grGXKEPqVUC1tx6G5Hw 5.16 BTC		
Key Size Key Usage Signature	1024 bits Encrypt, Verify, Derive 128 bytes : 84 FB EF B0 38 C7 B4 E6 BB 74 41 34 6C 52 6F 0E F7 7B 67 D4 D7 AD A7 EF 03 36 EF C0 40 DE 80 D3 0A A6 3C C9 C3 9E 31 52 07 EA 0F 44 63 EA 51 C2 D0 CE 9B A1 6D 99 F8 29 39 2C 40 98 67 12 8C A7 BF 43 30 D2 02 63 B3 F8 0C 36 85 50 8F 29 C2 87 CD E6 F0 EA 57 F2 A4 81 0E 94 B0 D8 46 67 2C 20 AE B6 89 64 A6 B0 DA 8E CF D4 09 7F 57 DC 9C 32 05 06 62 B2 08 14 FE 8B EA 11 13 BC F7 3B 7E F5		The server's host key is not cached in the registry. You have no guarantee that the server is the computer you think it is. The server's rsa2 key fingerprint is: ssh-rsa 1024 0a:27:d5:4f:00:9c:d1:a3:ff:ad:5c:cd:b3:7c:83:42 If you trust this host, hit Yes to add the key to PuTTY's cache and carry on connecting. If you want to carry on connecting just once, without adding the key to the cache, hit No. If you do not trust this host, hit Cancel to abandon the connection.	

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Standard model is useless against side channels and fault injection







RKA (related key attack) security considers these attacks



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Standard model is useless against side channels and fault injection







RKA (related key attack) security considers these attacks



Our research can help create better RKA schemes



















A wins if (1) Verify(pk,m',σ') =1 and (2) it didn't query m' to oracle



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A wins if (1) Verify(pk,m',σ') =1 and (2) it didn't query m' to oracle

Problem: This assumption is violated by tampering [AK96,...], side channels [W91,KJJ99,...] and fault injection [BS97,BdML97,...]

Attack on RSA-CRT [BdML97,L97] factors N given one faulty signature (attack also applies to Rabin signatures, and general RSA)
























A wins if (1) Verify(pk,m', σ') =1 and (2) it didn't get σ' from oracle (when querying on ϕ =id)



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 Φ ={id} gives standard unforgeability



A wins if (1) Verify(pk,m', σ') =1 and (2) it didn't get σ' from oracle (when querying on $\phi=id$) $\Phi=\{id\}$ gives standard unforgeability

 $\dot{\Phi}$ ={all functions} isn't possible [BK03]



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[BdML97] shows that RSA-CRT is not ϕ -RKA-secure for non-trivial ϕ



A wins if (1) Verify(pk,m', σ') =1 and (2) it didn't get σ' from oracle (when querying on ϕ =id)



[BdML97] shows that RSA-CRT is not *q*-RKA-secure for non-trivial *q*

Problem: φ-RKA schemes are really hard to construct (for interesting classes φ; for most primitives)

















Signature inherits **Φ-RKA** security from one-way function



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Many natural one-way functions (e.g., RSA) are Φ-RKA-secure with no additional assumptions



Signature inherits Φ-RKA security from one-way function

Many natural one-way functions (e.g., RSA) are Φ-RKA-secure with no additional assumptions

Creating new RKA-secure signatures is easier!





Φ-RKA-secure one-way functions are natural



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• RSA function is secure w.r.t. exponentiation



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Φ-RKA-secure one-way functions are natural

- RSA function is secure w.r.t. exponentiation
- Exponentiation (f(x)=g^x) is secure w.r.t. linear functions
- Learning with errors (f(s,e) = As+e mod q) is secure w.r.t. addition

Our result can pave the way for easier RKA constructions

My research



My research



What is Bitcoin?

What is Bitcoin?





What is Bitcoin?


















Centralized

Real-world identities Non-public transactions





Centralized

Real-world identities Non-public transactions



Public transactions



Centralized

Real-world identities Non-public transactions Regulated



Public transactions



Centralized

Real-world identities Non-public transactions Regulated



Pseudonyms Public transactions Unregulated*



Centralized

Real-world identities Non-public transactions Regulated Not anonymous



Pseudonyms Public transactions Unregulated*



Centralized

Real-world identities Non-public transactions Regulated Not anonymous



Pseudonyms Public transactions Unregulated* Potentially anonymous

(U) Bitcoin Virtual Currency: Unique Features Present Distinct Challenges for Deterring Illicit Activity

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How much anonymity does Bitcoin really provide?















1,200



1,200





























How do bitcoins get spent?

- decentralized
- transfer money
- generate money
- prevent double-spending
- decentralized
- transfer money
- generate money
- prevent double-spending



- decentralized
- transfer money
- generate money
- prevent double-spending



- decentralized
- transfer money
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To spend bitcoins, a user must indicate the previous transaction



- transfer money
- generate money
- prevent double-spending



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Collapse into a more manageable graph of clusters of public keys representing distinct entities

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Collapse into a more manageable graph of clusters of public keys representing distinct entities



Collect ground truth data by participating in transactions

























Heuristic #1: the same user controls these addresses [N08,RH11,RS13,A+13]





























In the standard idiom, change addresses are used at most twice: to receive change and to spend it

Clustering by change





Clustering by change





Clustering by change





Heuristic #2: the same user also controls this address





Engaged in transactions with:



Engaged in transactions with:

Exchanges





Engaged in transactions with:

• Exchanges



Vendors





Engaged in transactions with:

- Exchanges
- Mining pools



50 BTC

Vendors





Engaged in transactions with:

- Exchanges
- Mining pools



Vendors








Engaged in transactions with:

- Exchanges
- Mining pools



- **50 BTC**
- Vendors







Wallet services StrangCoin



Engaged in transactions with:





Engaged in transactions with:



Scraped published tags



Engaged in transactions with:



Scraped published tags

Found addresses discussed on forums





















Interacted with 31 MtGox addresses, tagged 518,723!

Participated in 344 transactions and tagged 1.3M public keys

Clustering using Heuristic 2



Clustering using Heuristic 2



Clustering using Heuristic 2











Can see when bitcoins meaningfully cross cluster boundaries



Identifying recipients potentially de-anonymizes user

Can see when bitcoins meaningfully cross cluster boundaries



Identifying recipients potentially de-anonymizes user

Hypothesis: if you subpoena exchanges, you can identify users















How much anonymity does Bitcoin really provide?



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Our analysis provides a real-world way to track flows of bitcoins



How much anonymity does Bitcoin really provide?

Our analysis provides a real-world way to track flows of bitcoins

Seems hard to launder significant quantities of money

My research


Acknowledgements





Acknowledgements



Thanks! Any questions?